

Cooperative management of IT and Electrical systems

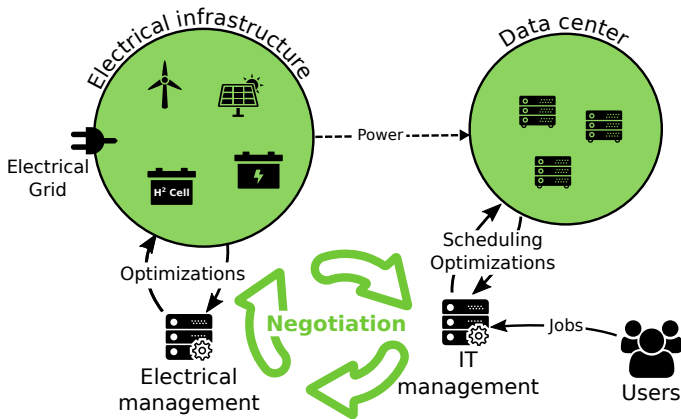
Georges Da Costa
Georges.Da-Costa@irit.fr

Workshop of renewable energy-powered datacenters
October 15th, 2019



Context

DATAZERO : the big picture



- ▶ Each DM has one or more objectives to satisfy
- ▶ Objectives may differ between DM
 - ▶ QoS related for ITDM, environmental impact for PDM

Managing different objectives

3 options studied :

- ▶ Finding a set of good solutions (set of possible trade-offs) (See previous presentation)
- ▶ Maximizing the weighted sum of the utilities, under the constraint of a distance between the two resulting profiles (SAN approach)
- ▶ Playing a game between the PDM and the ITDM so that each one maximizes its profit (GAN approach)

- ▶ Each DM has one or more objectives to satisfy
- ▶ Objectives may differ between DM
 - ▶ QoS related for ITDM, environmental impact for PDM

Managing different objectives

3 options studied :

- ▶ Finding a set of good solutions (set of possible trade-offs) (See previous presentation)
- ▶ Maximizing the weighted sum of the utilities, under the constraint of a distance between the two resulting profiles (SAN approach)
- ▶ Playing a game between the PDM and the ITDM so that each one maximizes its profit (GAN approach)

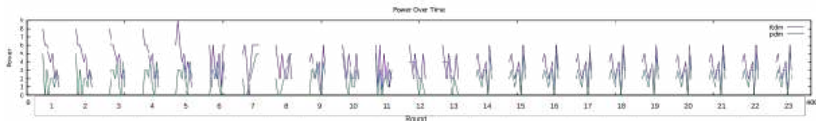
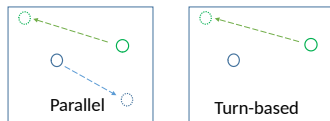
1. Context
2. SAN and GAN : Turn based approaches
3. SCHEDULING BASED NEGOTIATION (SAN)
Example
4. GAME BASED NEGOTIATION (GAN)
Model
Algorithm
Example
5. EXPERIMENT
6. CONCLUSION

SAN and GAN : Turn based approaches

Overview of SAN and GAN approaches

Main points

- ▶ Both algorithms are based on scheduling
 - ▶ DMs generate multiple scheduling solutions
 - ▶ Then we find negotiation solution from those scheduling solutions
- ▶ Both SAN and GAN negotiates in turn-based strategy
 - ▶ When ITDM runs scheduling (to follow PDM), PDM does not, and vice versa
 - ▶ We define 2 modes : "Follow PDM" mode (FLW_PD) and "Follow ITDM" mode (FLW_IT)
 - ▶ For both SAN and GAN, the whole system is executed **under only 1 mode** at a time



SCHEDULING BASED NEGOTIATION (SAN)

Definitions

- ▶ The set of ITDM profiles is $\{x_1, x_2, \dots, x_m\}$
- ▶ The set of PDM profiles is $\{y_1, y_2, \dots, y_n\}$
- ▶ Depending on each specific context, a profile may also be named "hint" or "candidate"

2 stages

- ▶ Stage 1 : Checking for matched pair
 - ▶ Decision variable : the pair {PDM profile, ITDM profile} :
 $\{x \in \{x_1, x_2, \dots, x_m\}, y \in \{y_1, y_2, \dots, y_n\}\}$
 - ▶ Objective : maximize sum of utility

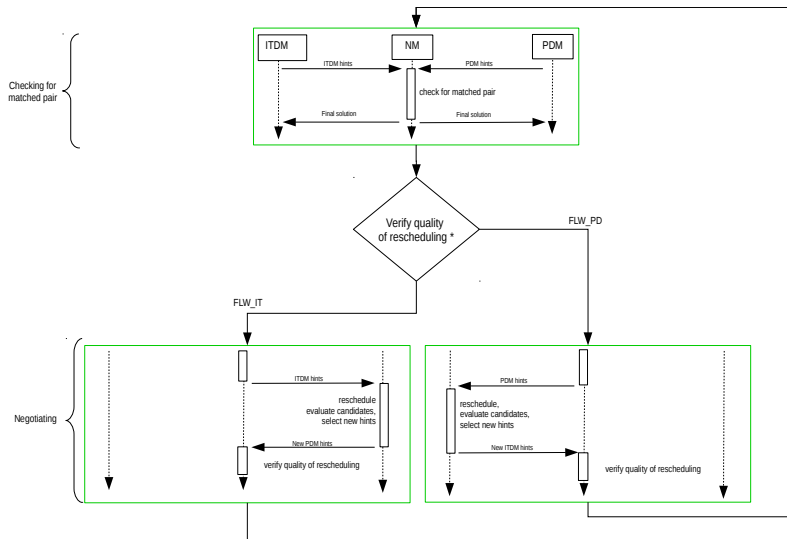
$$\max_{\{x,y\}} (u(x) + u(y)) \quad (1)$$

- ▶ Constraint : $d(x, y) < \epsilon$
where $d()$ is the distance between x and y
- ▶ If we can't find any matched pair, run Stage 2 : Negotiating.

Stage 2 : Negotiating :

- ▶ General mechanism
 - ▶ At a time, the whole system is executed **under only 1 mode** : follow ITDM (FLW_IT) or follow PDM (FLW_PD)
 - ▶ NM decides to switch between two modes using "*verify quality of rescheduling*"
 - ▶ Repeat until matched pair found
- ▶ Two modes :
 - ▶ FLW_IT : Follow the ITDM
 - ▶ NM sends the ITDM hints to PDM
 - ▶ PDM uses an algorithm (e.g. greedy, linear program) to find multiple scheduling solutions as candidates
 - ▶ PDM evaluates quality of candidates by "*weighted similarity*" to hints
 - ▶ PDM selects candidates with high "*weighted similarity*" as its news hints and sends back to NM
 - ▶ FLW_PD : Similar to FLW_IT, following the PDM

Algorithm



(*) verify quality of rescheduling: compare "distance between the best ITDM hint and the best PDM hint" before and after rescheduling

Example

Example

Follow ITDM



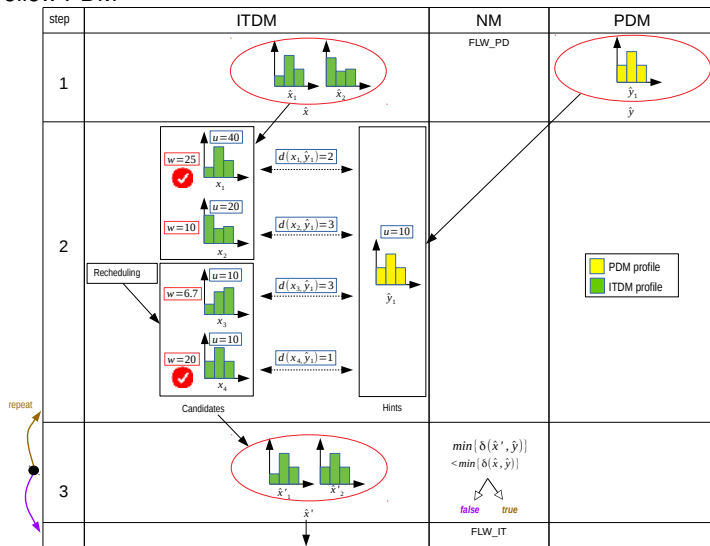
Notation:
 - candidates don't have "hat"
 - hints have "hat"
 - hints of next round have "prime"

$$w = \frac{4\hat{x}_1 + 15}{d(\hat{x}_2, y_1)} + \frac{\hat{x}_1 + 15}{d(\hat{x}_2, y_1)} = \frac{61}{5} + \frac{21}{2} = 12 + 10 = 22$$

step	ITDM	NM	PDM
1		FLW_IT	
2		<p>Hints</p>	<p>Candidates</p> <p>Recheduling</p>
3	<p>repeat</p>	<p>Evaluate</p> $\min\{\delta(\hat{x}, \hat{y}^*)\}$ $< \min\{\delta(\hat{x}, \hat{y})\}$ <p>false true</p>	
		FLW_PD	
		...	

Example

Follow PDM





GAME BASED NEGOTIATION (GAN)

Model



Game players

ITDM:

- Objective:  
 - max {payment from users – payment to PDM}
- Decision variables:
 - price
 - scheduling

revenue - cost

PDM:

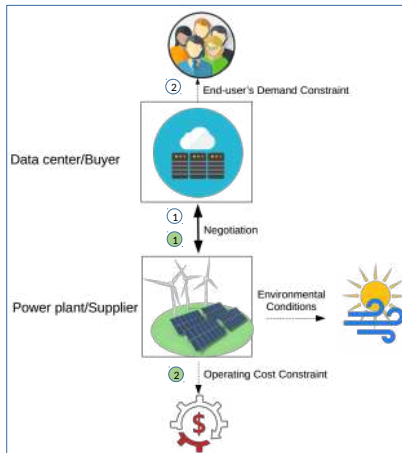
- Objective:  
 - max {payment from ITDM – (opex + capex cost)}
- Decision variables:
 - purchased power
 - scheduling

Game model

Hybrid model

- Non-cooperative: each player maximizes their own utility
- Cooperative: sometimes a player follows the other's suggestion

Supplier-buyer game diagram



Motivation

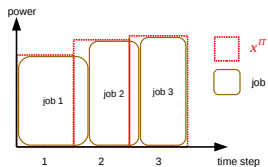
- ▶ The buying-selling nature of the system
 - ▶ Selling : PDM is selling power → controls the price
 - ▶ Buying : ITDM is buying power → decide the order/purchase
- ▶ Advantages of pricing
 - ▶ Power source availability can be reflected in price
 - ▶ Through price, pattern of order reflects pattern of PDM's desirable supply → drive demand toward supply

Preliminary

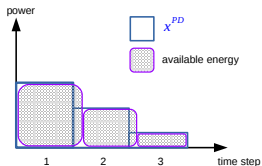
- ▶ Players are selfish
 - ▶ They try to maximize their utility
 - ▶ Each player negotiates just because he foresees some benefit
 - ▶ We introduce *incentive pricing mechanism* : each player tries to find offers that are attractive to the other player.
- ▶ An unexpected situation may occurred : all players can't foresee their benefit and stop negotiate without reaching any agreement.
 - ▶ From the view of the whole system, this situation is unacceptable, no transaction is done, the players obtain zero utility
 - ▶ If this situation occurred, we introduce *sacrifice mechanism*, in which the players gradually sacrifice their utility until they reach an agreement.

Definition 1

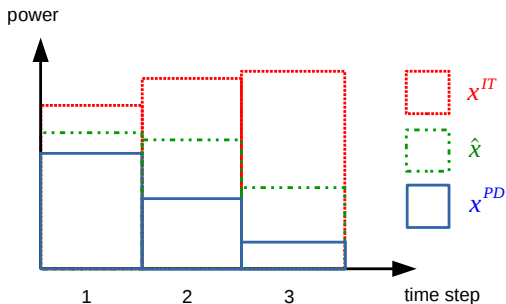
- ▶ T : Time window
- ▶ \hat{x}, x^{IT}, x^{PD} are profiles
- ▶ $\hat{X} = \{\hat{x}_1, \hat{x}_2, \dots, \hat{x}_T\}$, $X^{IT} = \{x_1^{IT}, x_2^{IT}, \dots, x_T^{IT}\}$ $X^{PD} = \{x_1^{PD}, x_2^{PD}, \dots, x_T^{PD}\}$



(a)

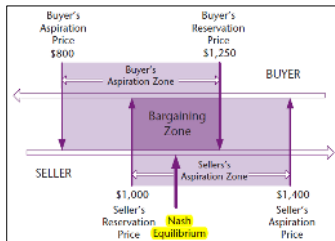


(b)

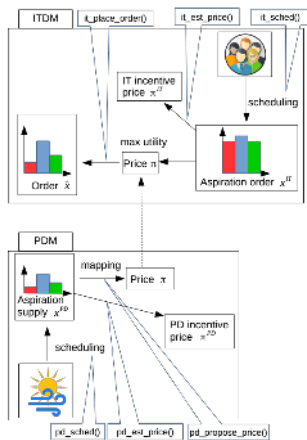


Definition 2

► $\pi = \{\pi_1, \pi_2, \dots, \pi_T\}$, $\pi^{IT} = \{\pi_1^{IT}, \pi_2^{IT}, \dots, \pi_T^{IT}\}$, $\pi^{PD} = \{\pi_1^{PD}, \pi_2^{PD}, \dots, \pi_T^{PD}\}$



2006_Krause_Bargaining Stances and Outcomes in Buyer-Seller Negotiations: Experimental Results (The Journal of Supply Chain Management)

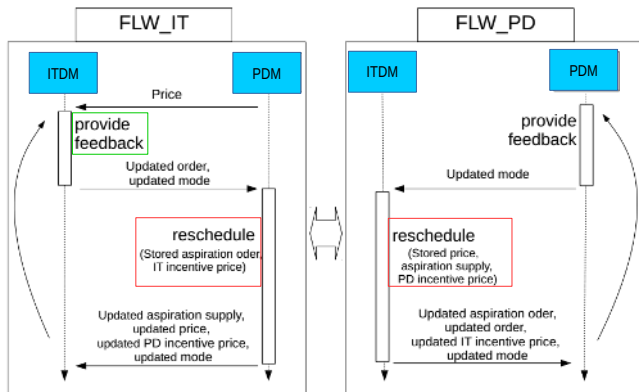


Algorithm

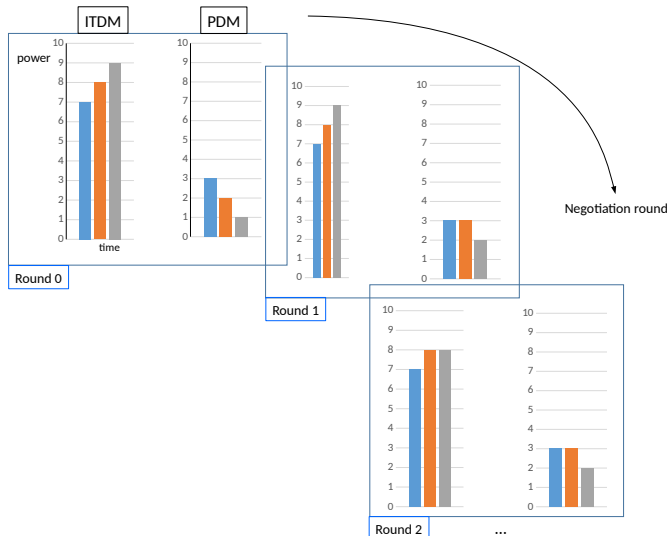
Algorithm

Note :

- ▶ At a time, the whole system is executed under **only one mode** : **FLW_IT** or **FLW_PD**
- ▶ DMs only exchange data that have been updated/modified, other data can be stored and reused



Graphical interpretation of the algorithm



Example

Example

Simplified example

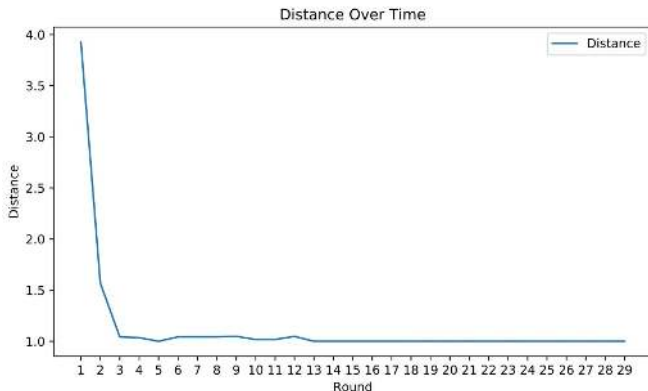
round	row	PD-Player	Common Space	IT-Player
1	1	<p>power</p> <p>time</p> <p>x^{PD}</p>	<p>$\pi^{PD} = [7,6]$</p> <p>$\pi = [8,6]$</p> <p>\hat{x}</p> <p>$\pi^{IT} = [3,3]$</p>	<p>x^{IT}</p> <p>Scheduling solution</p>
	2	<p>Assume $\begin{cases} r(x^{IT}, \pi^{IT}) + \alpha = 0.25 \\ r(\hat{x}, \pi) = 0.22 \end{cases}$</p> <p>$r(x^{IT}, \pi^{IT}) + \alpha > r(\hat{x}, \pi) \rightarrow \text{true}$</p>		<p>Assume $\begin{cases} c(x^{PD}, \pi^{PD}) - \alpha = 0.13 \\ c(\hat{x}, \pi) = 0.15 \end{cases}$</p> <p>$c(x^{PD}, \pi^{PD}) - \alpha < c(\hat{x}, \pi) \rightarrow \text{true}$</p>
	3	$pd_pre = FLW_IT \leftarrow$	$mod = FLW_IT$	$it_pre = FLW_PD \leftarrow$
2	4	<p>x^{PD}</p>	<p>$\pi^{PD} = [7,6]$</p> <p>$\pi = [7,6]$</p> <p>\hat{x}</p> <p>$\pi^{IT} = [3,3]$</p>	
	5	<p>Assume $\begin{cases} r(x^{IT}, \pi^{IT}) + \alpha = 0.25 \\ r(\hat{x}, \pi) = 0.28 \end{cases}$</p> <p>$r(x^{IT}, \pi^{IT}) + \alpha > r(\hat{x}, \pi) \rightarrow \text{false}$</p>		<p>Assume $\begin{cases} c(x^{PD}, \pi^{PD}) - \alpha = 0.16 \\ c(\hat{x}, \pi) = 0.17 \end{cases}$</p> <p>$c(x^{PD}, \pi^{PD}) - \alpha < c(\hat{x}, \pi) \rightarrow \text{true}$</p>
	6	$pd_pre = FLW_PD \leftarrow$	$mod = FLW_PD$	$it_pre = FLW_PD \leftarrow$
3	7		<p>$\pi^{PD} = [7,6]$</p> <p>$\pi = [7,6]$</p> <p>\hat{x}</p> <p>$\pi^{IT} = [6,5]$</p>	<p>x^{IT}</p> <p>New scheduling solution</p>
	8	<p>Assume $\begin{cases} r(x^{IT}, \pi^{IT}) + \alpha = 0.23 \\ r(\hat{x}, \pi) = 0.24 \end{cases}$</p> <p>$r(x^{IT}, \pi^{IT}) + \alpha > r(\hat{x}, \pi) \rightarrow \text{false}$</p>		<p>Assume $\begin{cases} c(x^{PD}, \pi^{PD}) - \alpha = 0.16 \\ c(\hat{x}, \pi) = 0.14 \end{cases}$</p> <p>$c(x^{PD}, \pi^{PD}) - \alpha < c(\hat{x}, \pi) \rightarrow \text{false}$</p>
	9	$pd_pre = FLW_PD \leftarrow$		$it_pre = FLW_IT \leftarrow$
	10	<p>$d(x^{PD}, \hat{x}) > \epsilon \rightarrow \text{true}$</p> <p>$d(x^{PD}, \hat{x}) > \epsilon \rightarrow \text{false}$</p>		<p>$d(x^{PD}, \hat{x}) > \epsilon \rightarrow \text{true}$</p> <p>$d(x^{PD}, \hat{x}) > \epsilon \rightarrow \text{false}$</p>

EXPERIMENT

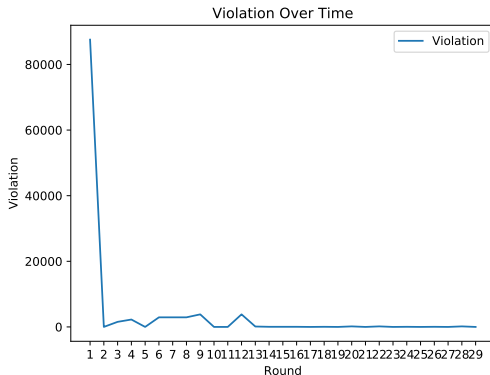
Real PDM & ITDM

- ▶ PDM weather information : 1 month
- ▶ Time window : 3 days or 72 hours
- ▶ Timestep : 1 hour or 3,600,000 ms
- ▶ PDM sizing : $\approx 1\text{kW}$
- ▶ Run time : ≈ 10 minutes

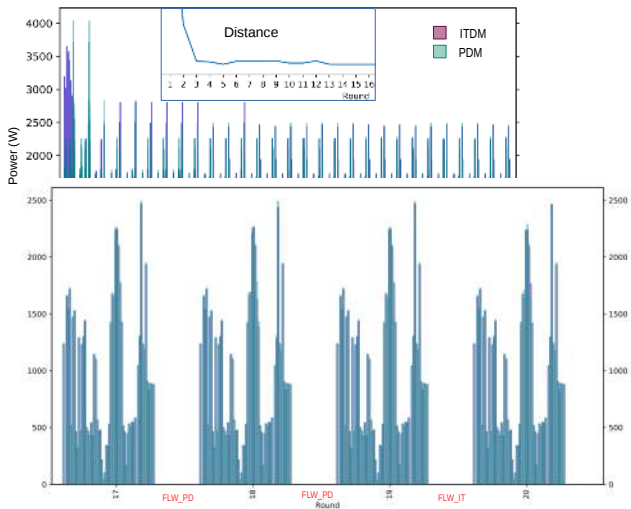
- ▶ Calculation : using Pearson correlation
- ▶ Distance is not always decreasing because the profiles are evaluated by both utility and distance
- ▶ Negotiation results depend a lot on the series of utilities from DMs



- ▶ Calculation : sum of the amount that the ITDM profiles exceeds PDM's profiles
- ▶ A significant reason for this result : DMs scheduling algorithms



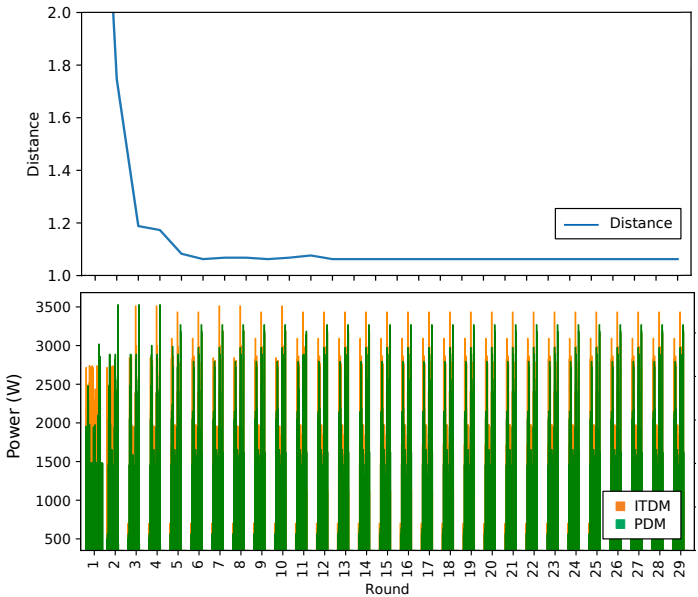
Power level



Real PDM & ITDM

- ▶ PDM weather information : 1 month
- ▶ Time window : 3 days or 72 hours
- ▶ Timestep : 1 hour or 3,600,000 ms
- ▶ PDM sizing : $\approx 4\text{kW}$
- ▶ Run time : ≈ 10 minutes

Power level and convergence



CONCLUSION

- ▶ GAN and SAN provide (semi) black-box based negotiation for renewable-based datacenters
- ▶ Guaranty of convergence
- ▶ Semi black-box approach achieves higher stability and performance than the black-box one
- ▶ Perspectives
 - ▶ Performance evaluation of the algorithm itself (number of call to black boxes)
 - ▶ Multi-time scale decision
 - ▶ Comparison with GreenSlot

See you on datazero.org