

University of Stuttgart

Institute of Parallel and Distributed Systems -
Distributed Systems Group

$s, t) \leq \max$
abled intervals
We need $[-]^+$ in bs
 $(t, td, n)) - \beta_s(\max(s, te, n))$
interval $(s, t]$ starts in the ena
interval ends before t (cf. enabled
2), we have $R(t) - A(s) \geq \beta_s$



Modeling Time-Triggered Service Intermittence In Network Calculus

Jonathan Falk, Frank Dürr, Kurt Rothermel

The age of the cyber-physical machine.

More distributed systems interfacing with the physical world

- “Smart” {city, factory, home}
- Autonomous Driving
- IEEE Time-sensitive Networking (TSN) Workgroup

- Increasingly complex network setups and network behavior
- Often non-trivial analysis

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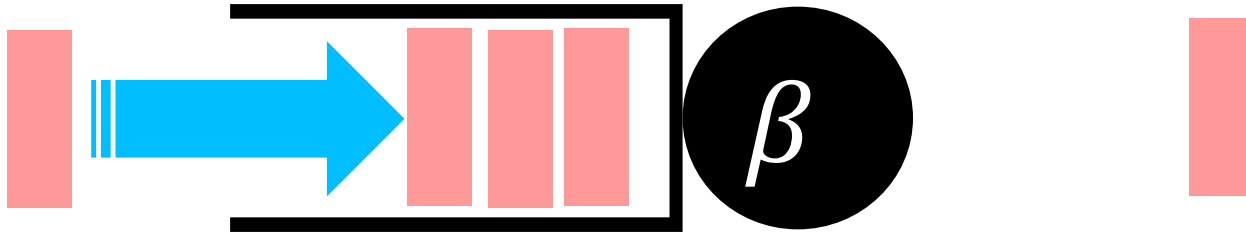
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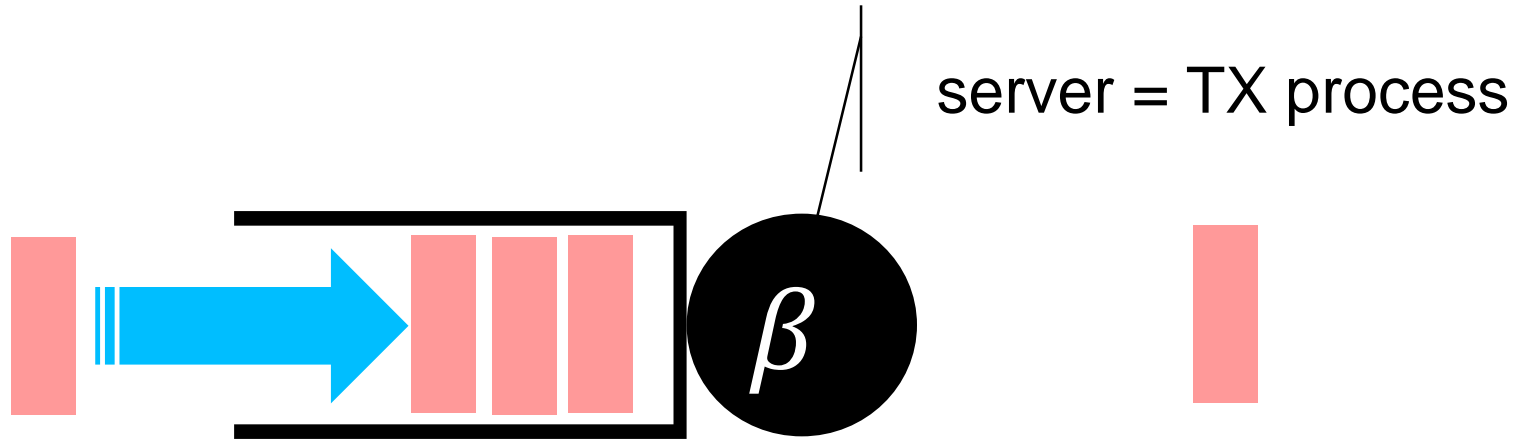


Formal frameworks

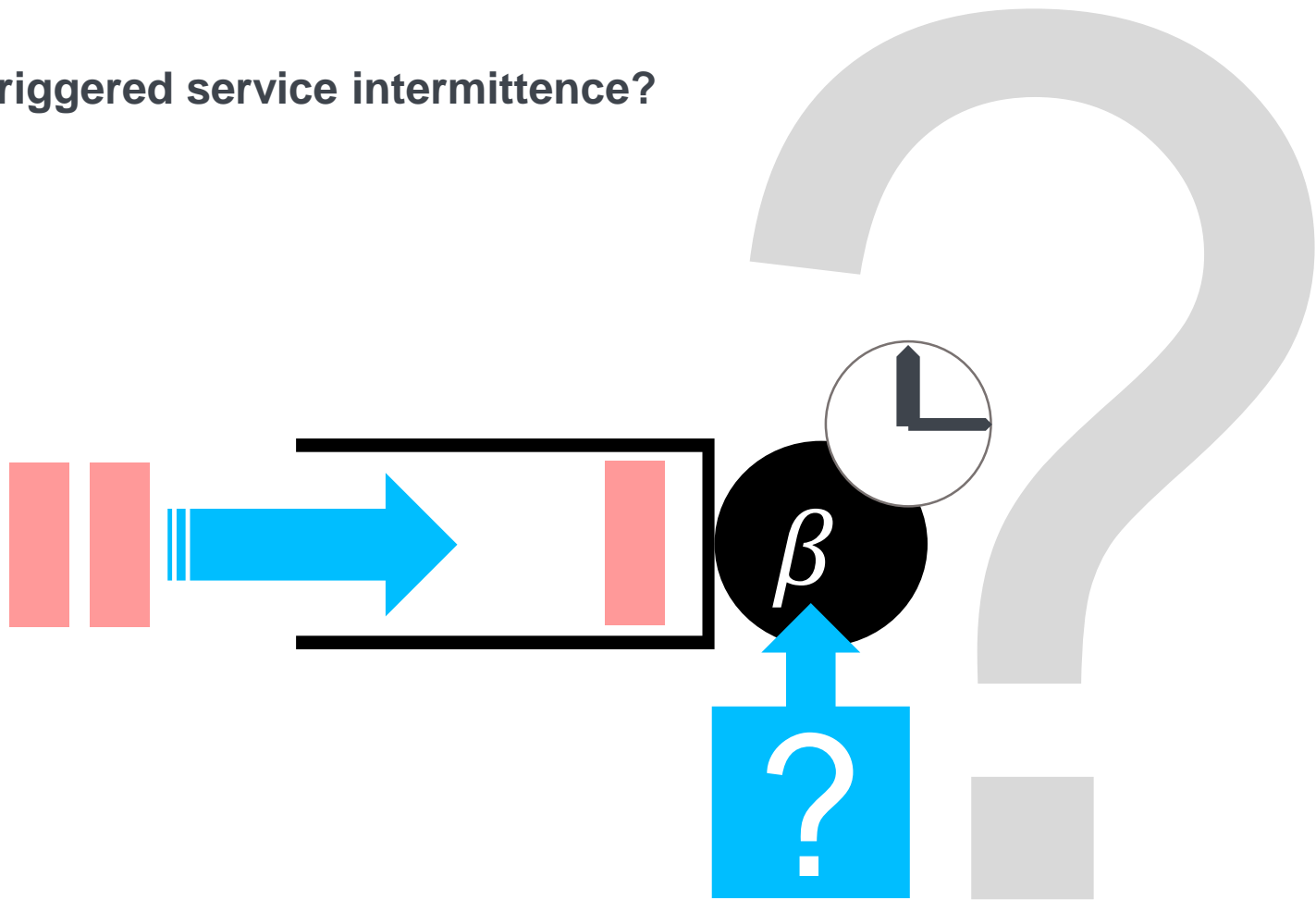
(Deterministic) network calculus to compute bounds in queuing systems



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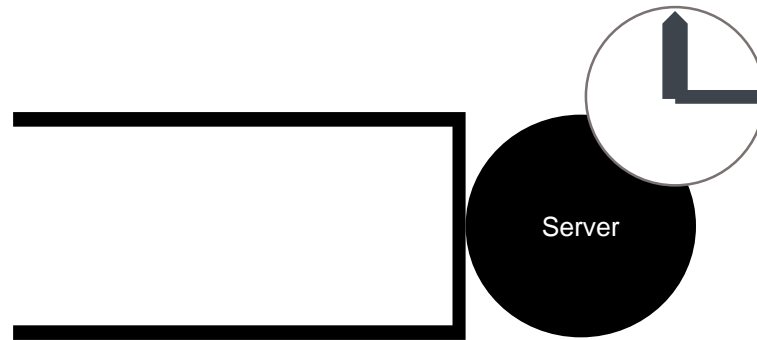
Time-triggered service intermittence?



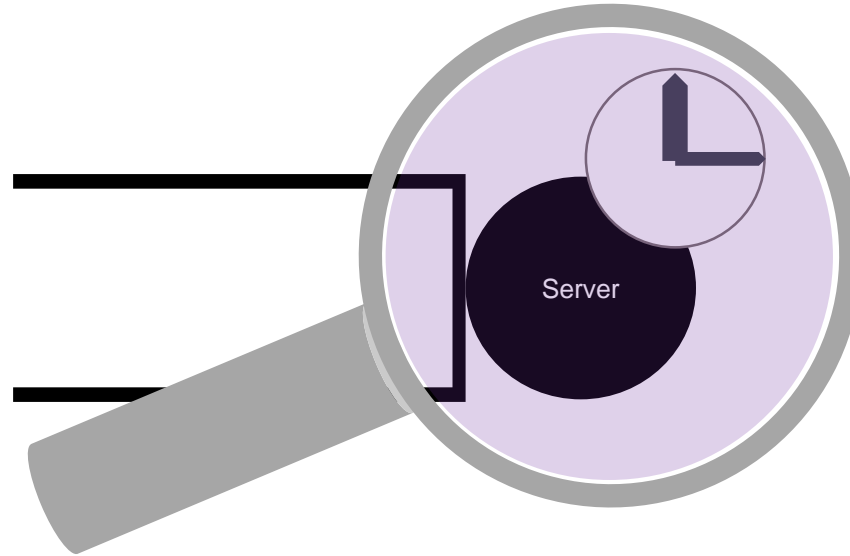
Contributions

- Modeling 2 classes of systems with time-triggered service intermittence
 - Time-variant
 - Time-invariant
 - Implications?
- Highly generic model
 - Complex schedules for service intermittence
 - Beyond constant-rate service

Systems with intermittent service



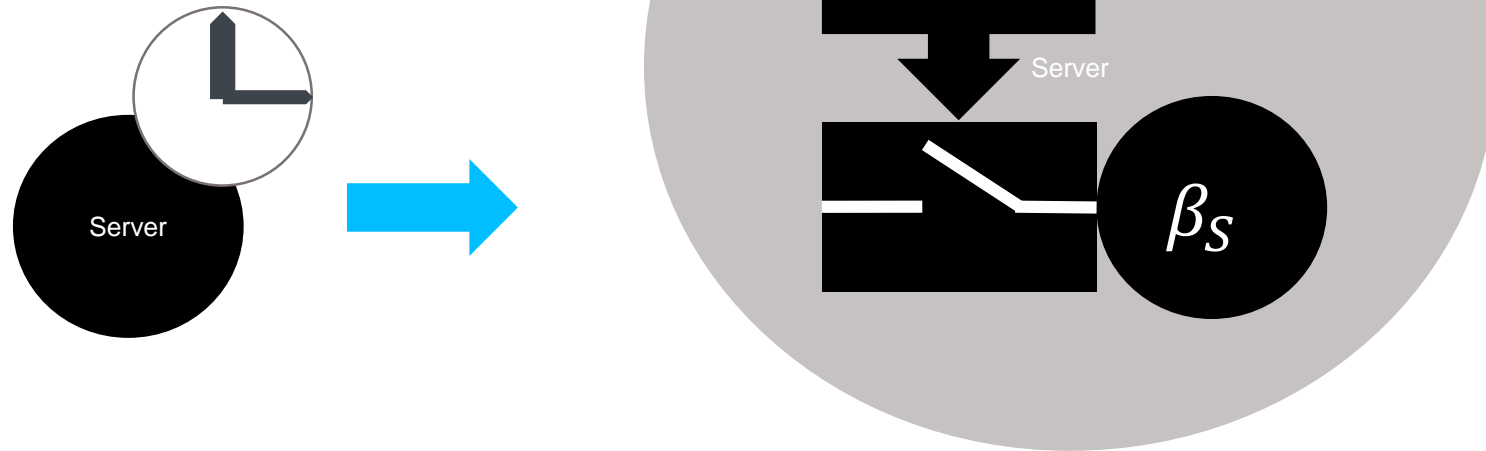
Systems with intermittent service



Time-triggered blocking

Motivating example:

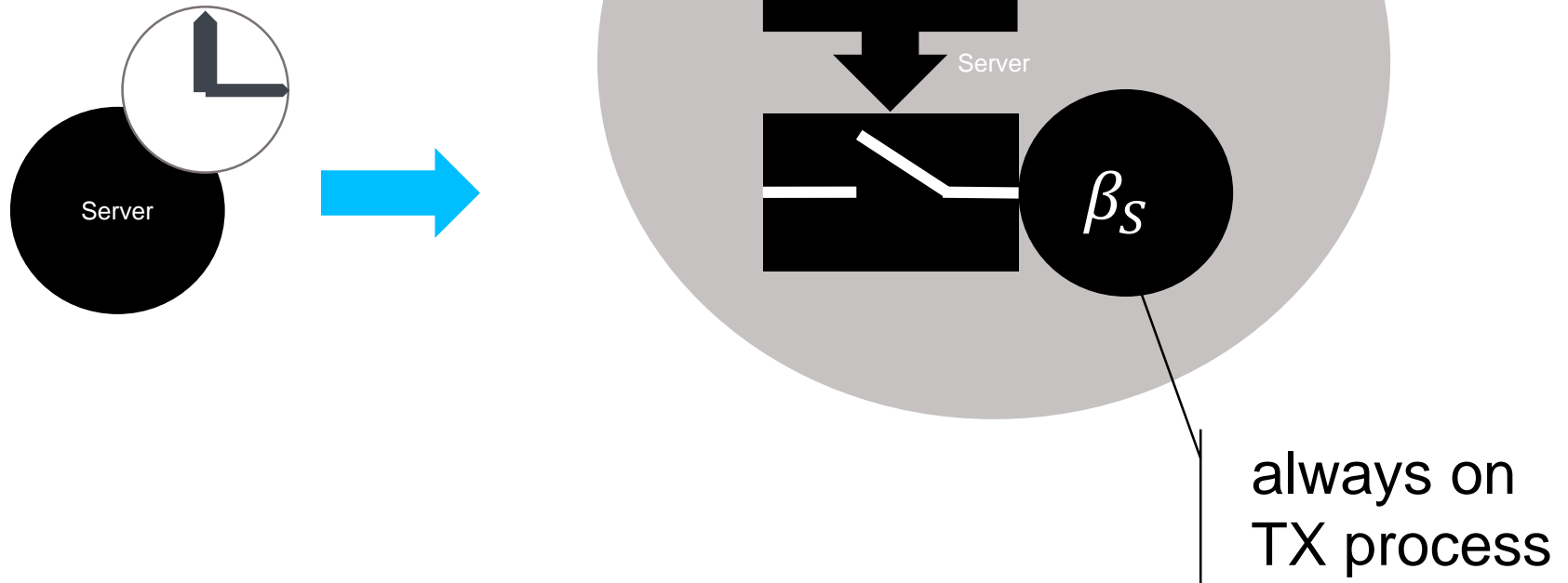
Time-aware shaping



Time-triggered blocking

Motivating example:

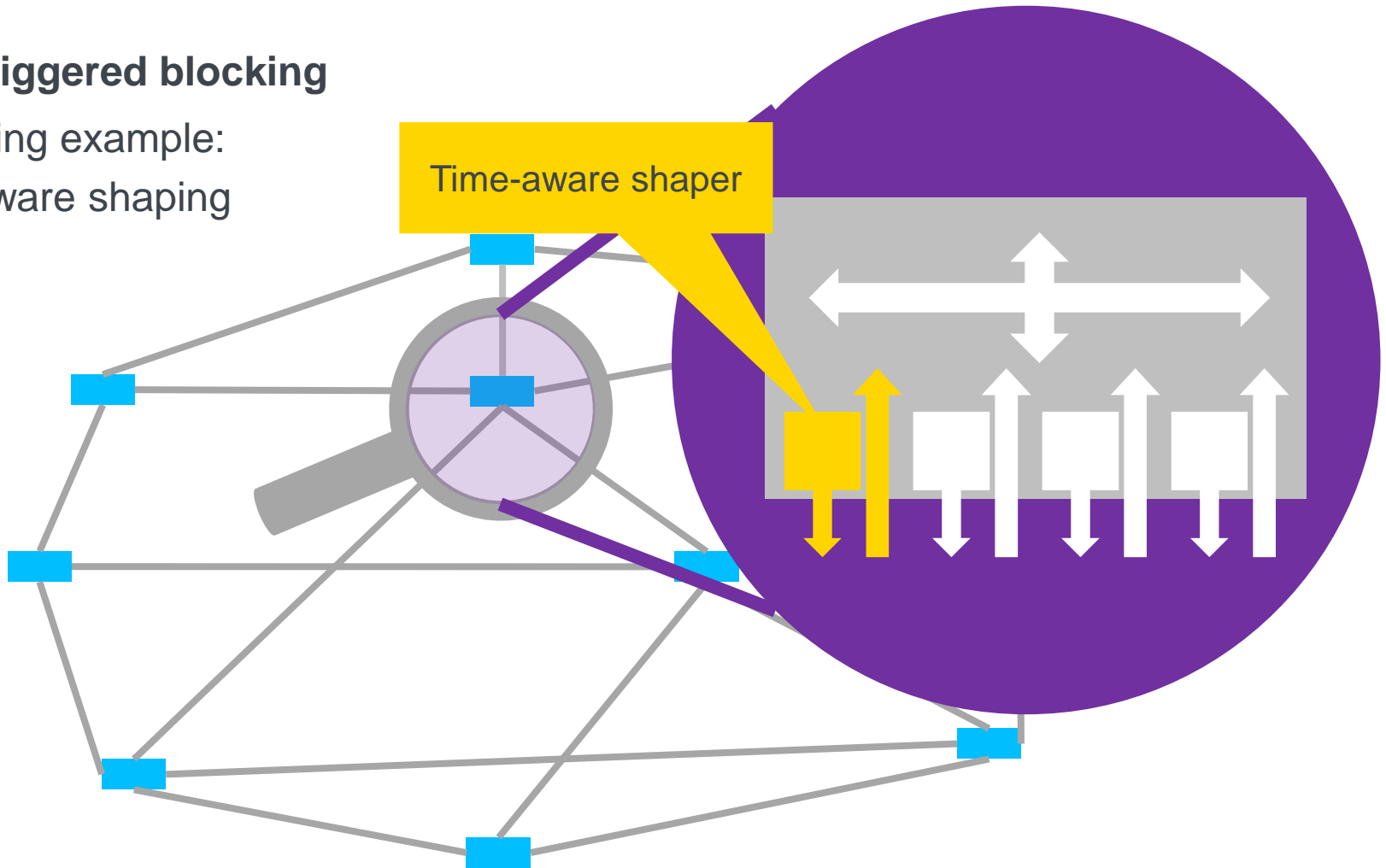
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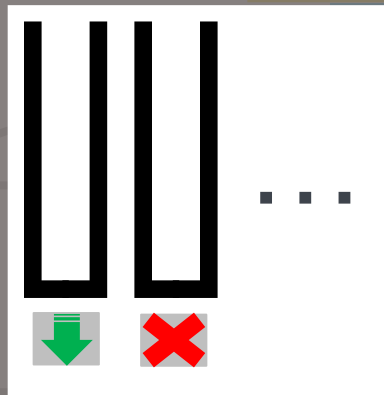
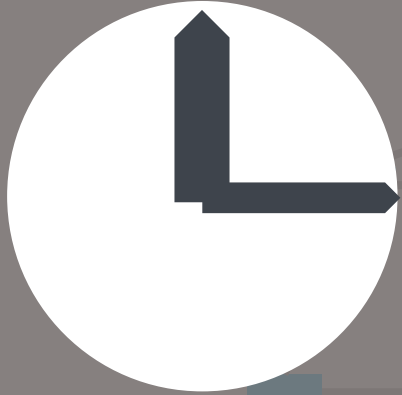
Time-aware shaping



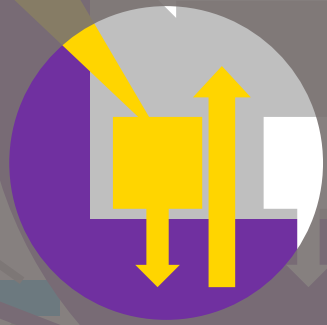
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Motivating example:

Time-aware shaping



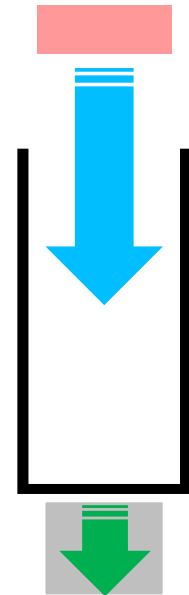
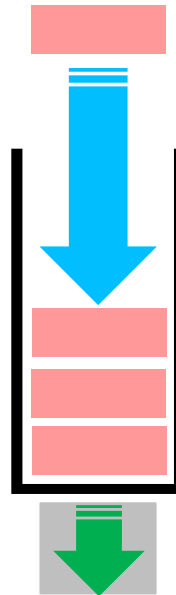
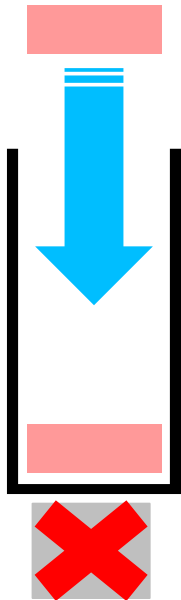
Time-aware shaper



Time	State
t_1	↓ ↓ ↓ ×
t_2	× × ↓ ↓ ↓
t_3	× × × ×

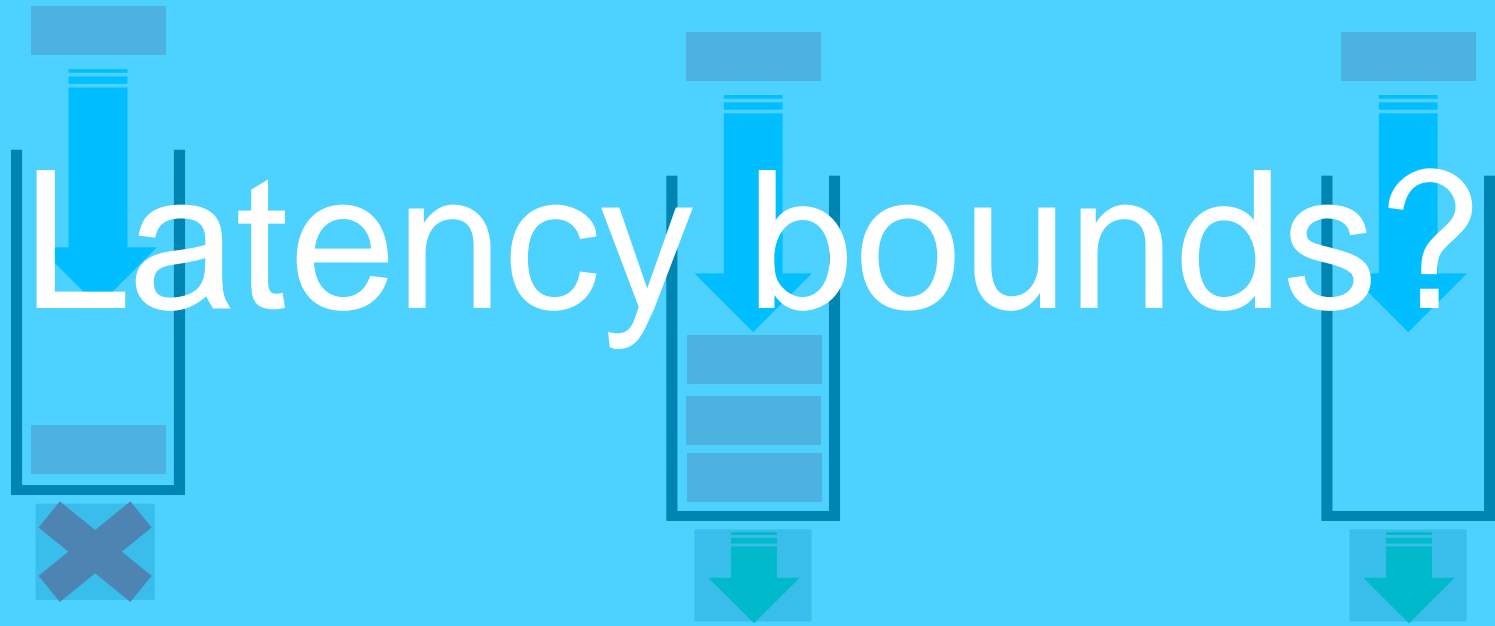
Time-triggered blocking

Data transmission and network elements are **not synchronized**.



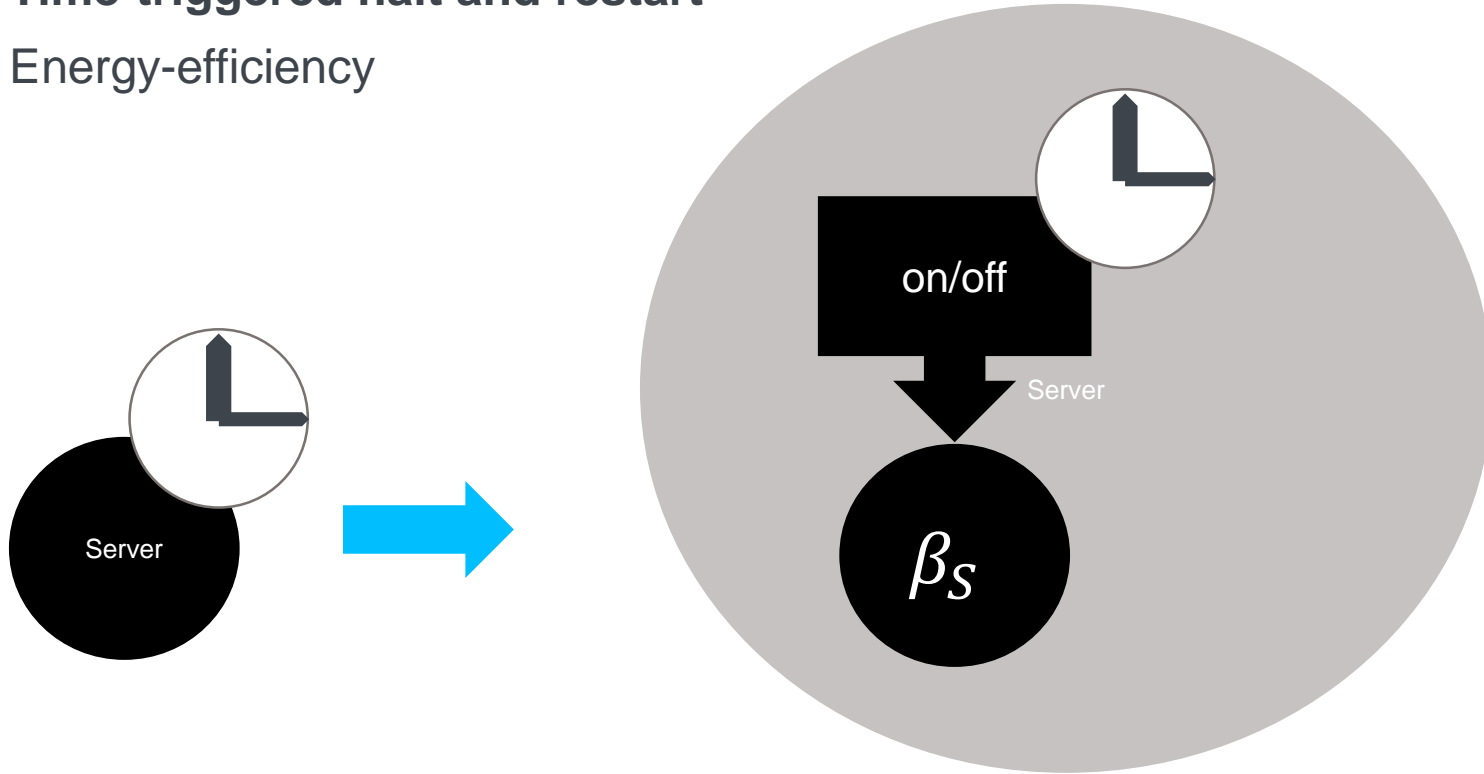
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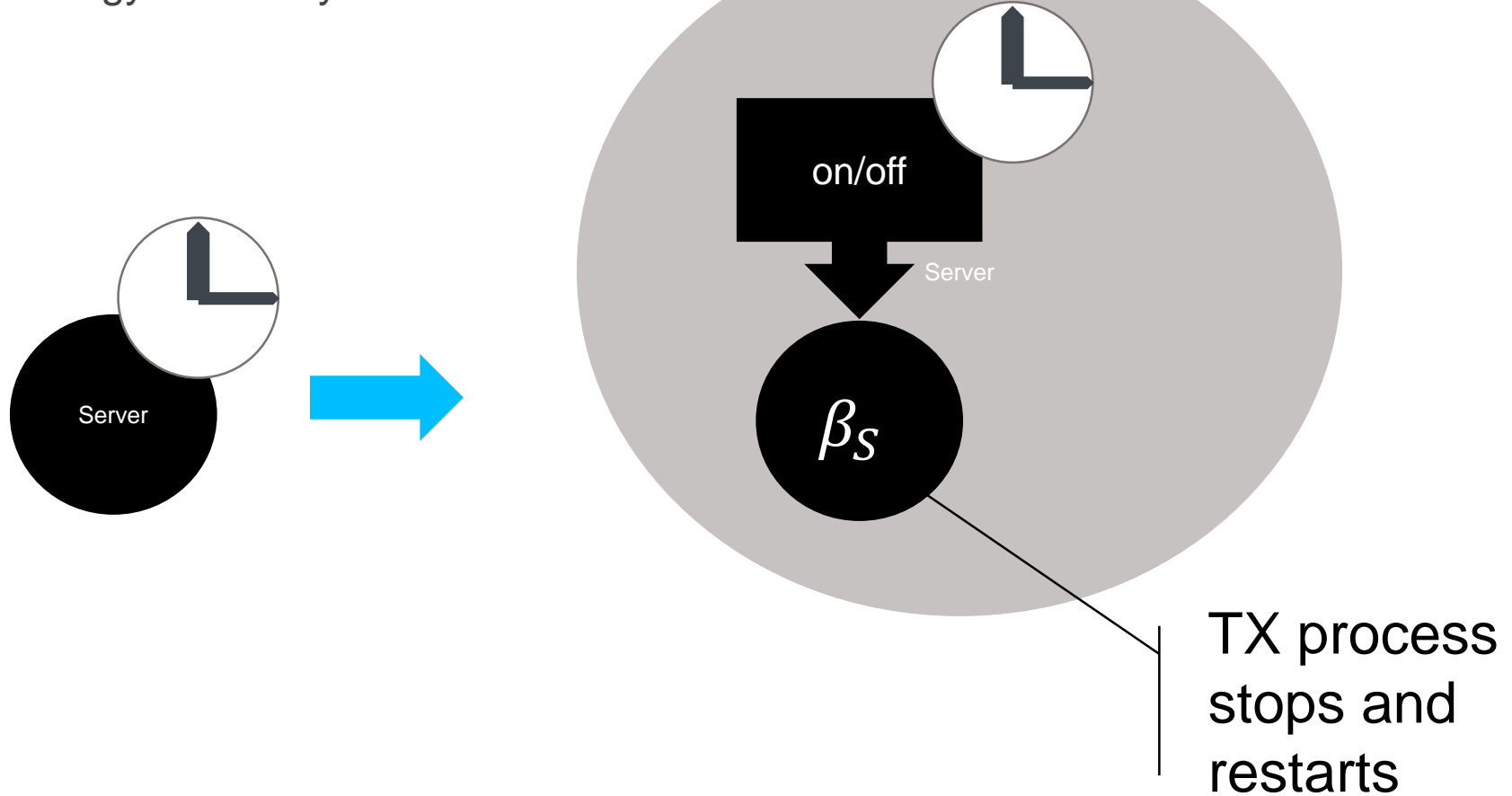
Time-triggered halt and restart

Energy-efficiency



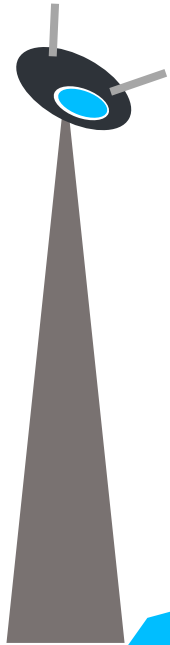
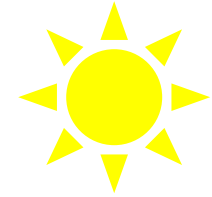
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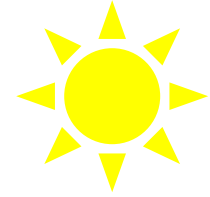
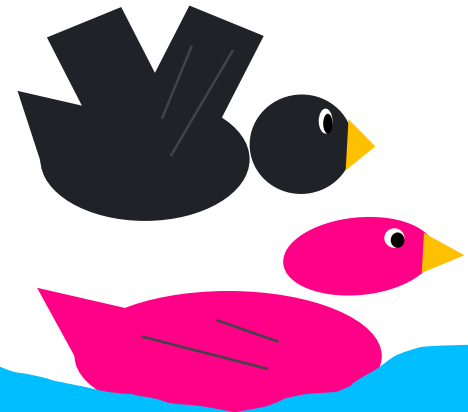
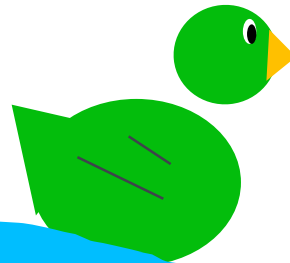
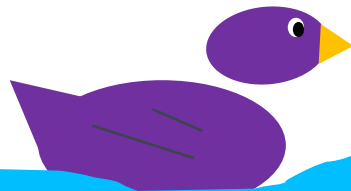
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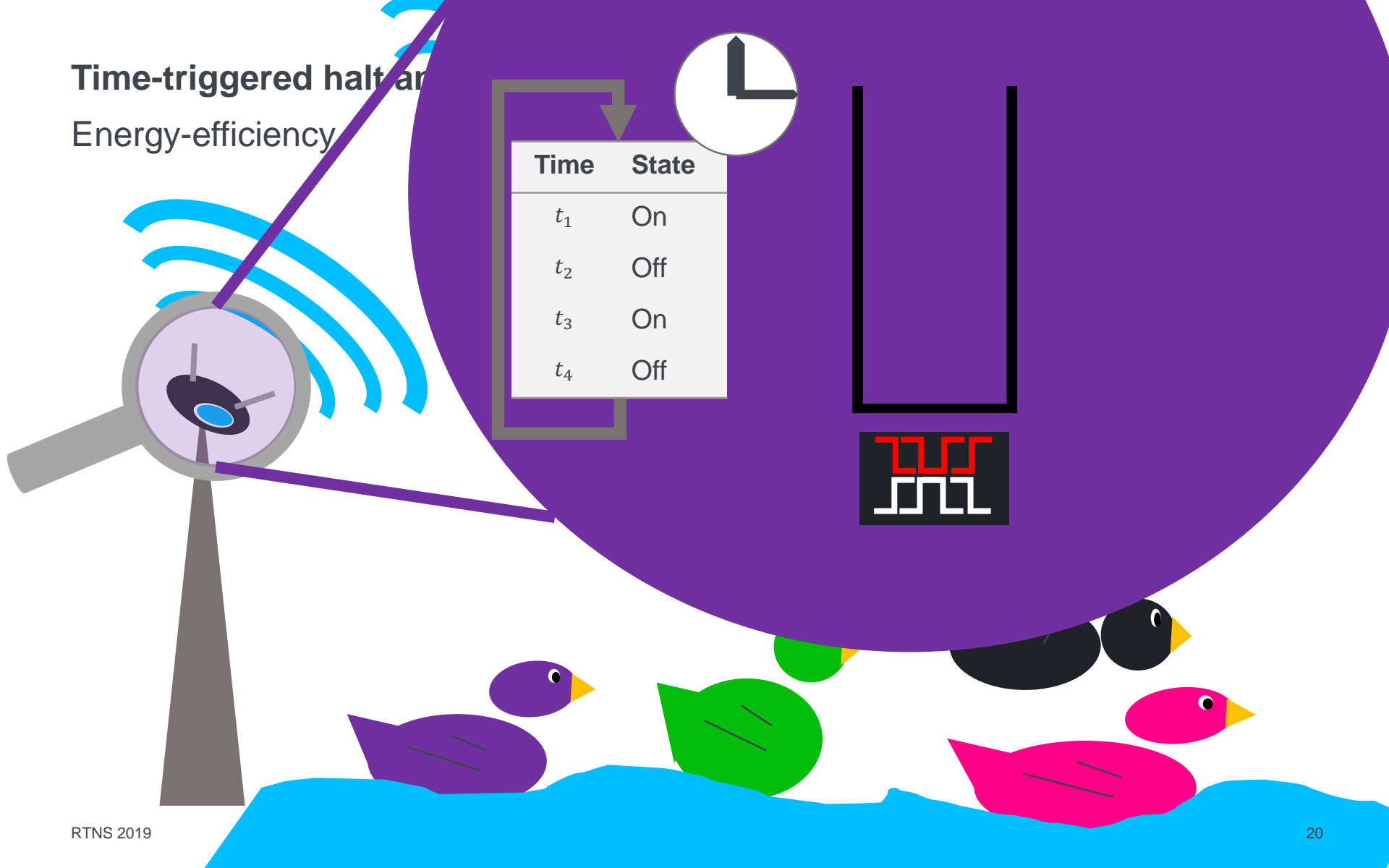


Time-triggered halt and restart

Energy-efficiency

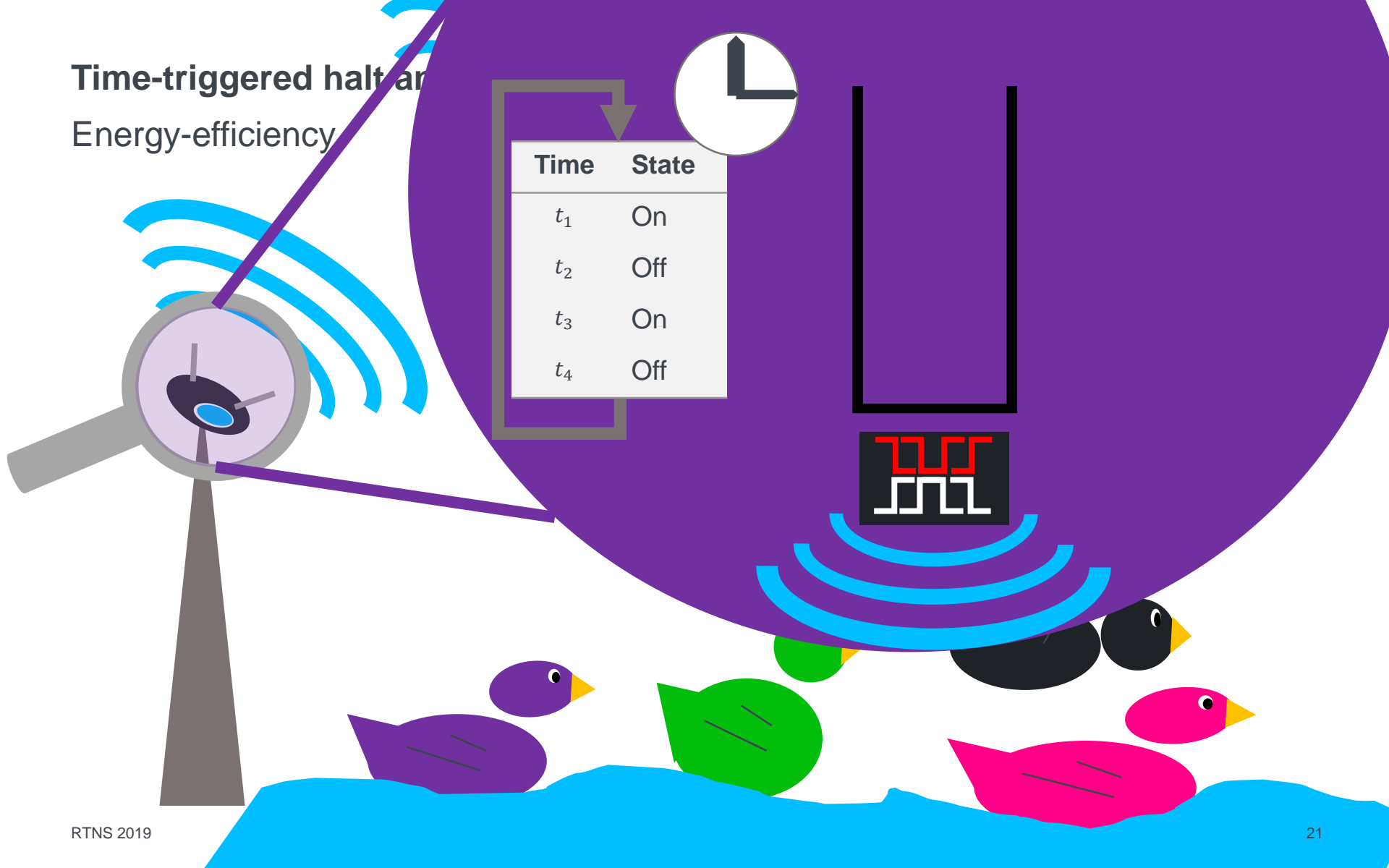


Time-triggered half an
Energy-efficiency



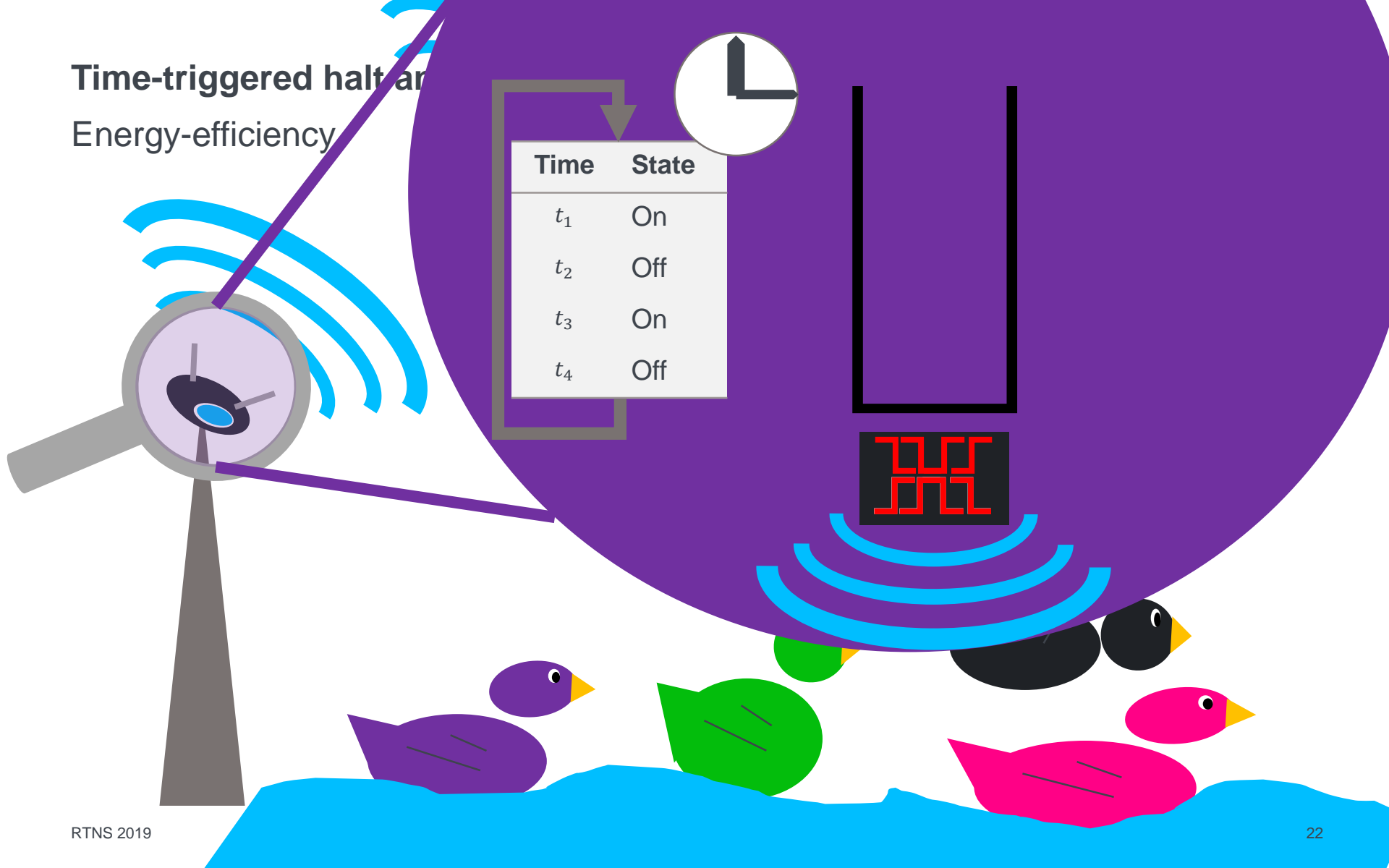
Time	State
t_1	On
t_2	Off
t_3	On
t_4	Off

Time-triggered half an
Energy-efficiency



Time	State
t_1	On
t_2	Off
t_3	On
t_4	Off

Time-triggered half an
Energy-efficiency



Time	State
t_1	On
t_2	Off
t_3	On
t_4	Off

Time-triggered half an
Energy-efficiency



A clock icon is positioned above the table. A diagram shows a box on the right with an arrow pointing to the table, and another arrow pointing from the table back to the box, forming a loop.

Time	State
t_1	On
t_2	Off
t_3	On
t_4	Off

Backlog Bounds?

Related work

- TDMA

- Gollan, N., and J. Schmitt. “Energy-Efficient TDMA Design Under Real-Time Constraints in Wireless Sensor Networks.” In Proceedings of the 2007 15th International Symposium on Modeling, Analysis, and Simulation of Computer and Telecommunication Systems, 80–87. MASCOTS '07. Washington, DC, USA.
- Dang Dinh Khanh, and Ahlem Mifdaoui. “Timing Analysis of TDMA-Based Networks Using Network Calculus and Integer Linear Programming.” In 2014 IEEE 22nd International Symposium on Modelling, Analysis Simulation of Computer and Telecommunication Systems (MASCOTS), 21–30. Paris, France.

- Ethernet

- Zhao, Luxi, Paul Pop, Qiao Li, Junyan Chen, and Huagang Xiong. “Timing Analysis of Rate-Constrained Traffic in TTEthernet Using Network Calculus.” *Real-Time Systems* 53, no. 2 (March 1, 2017): 254–87.
- Zhao, L., P. Pop, and S. S. Craciunas. “Worst-Case Latency Analysis for IEEE 802.1Qbv Time Sensitive Networks Using Network Calculus.” *IEEE Access* 6 (2018): 41803–15.
- Zhao, L., P. Pop, Z. Zheng, and Q. Li. “Timing Analysis of AVB Traffic in TSN Networks Using Network Calculus.” In 2018 IEEE Real-Time and Embedded Technology and Applications Symposium (RTAS), 25–36, 2018.

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- Dang Dinh Khanh, and Ahlem Mifdaoui. “Timing Analysis of Real-Time Systems Using Linear Programming.” In 2014 IEEE 22nd International Symposium on Real-Time Systems (MASCOTS), 21–30. Paris, France.

- Ethernet

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simple schedules

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- Ethernet

simple schedules

limited scope

Analysis of Rate-Constrained Traffic in TTEthernet Using Network Calculus.” In 2014 IEEE 22nd International Symposium on Modeling, Analysis and Simulation of Computer and Telecommunication Systems (MASCOTS), 21–30. Paris, France.

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- Ethernet

simple schedules

time-invariant

limited scope

Related work

ultimately constant-rate service curves*

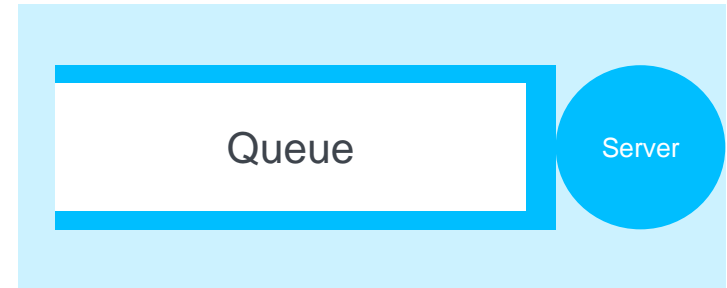
simple schedules

time-invariant

limited scope

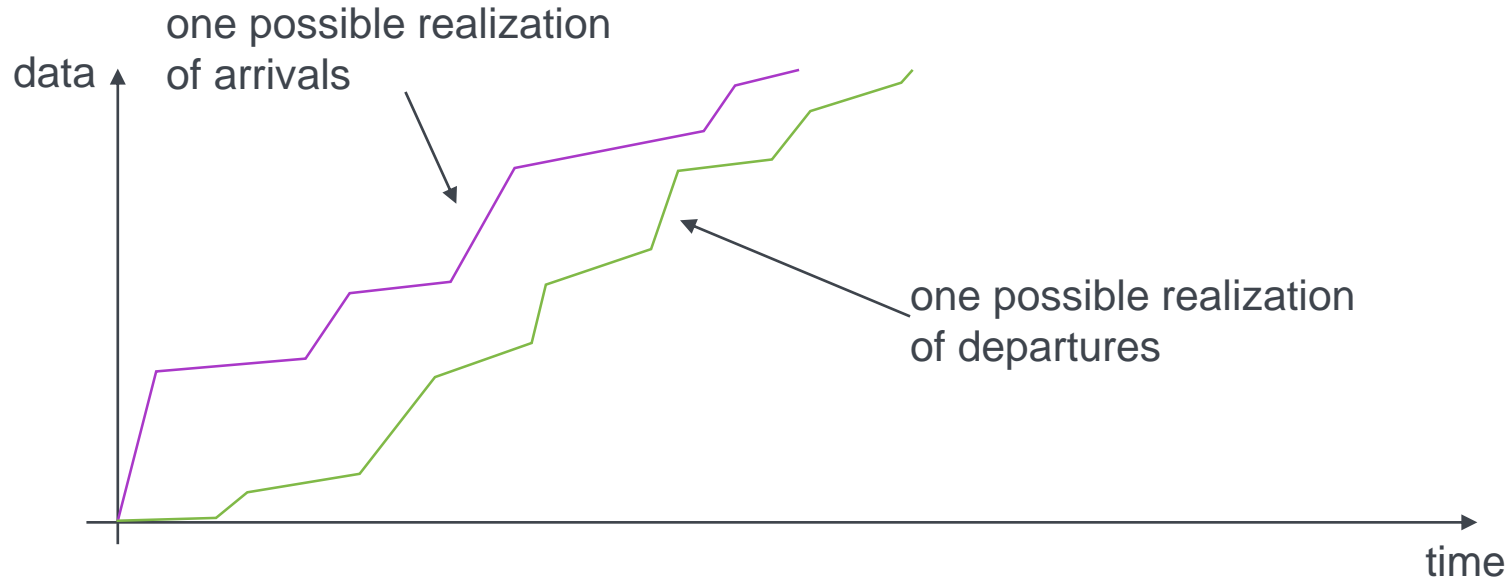
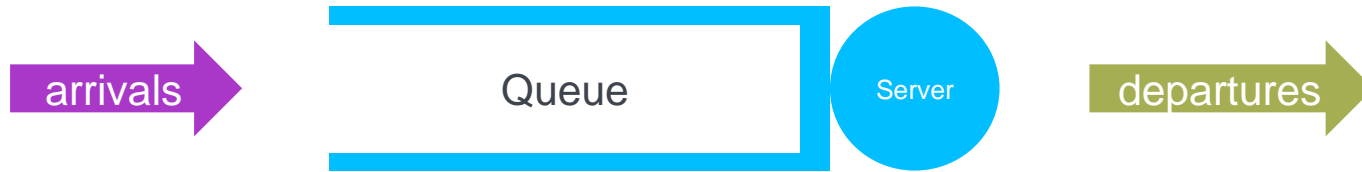
Network calculus

- “System theory” for queuing systems
 - Modularity (cf. convolution operation)
 - Deterministic (guaranteed) bounds
- NC allows to compute **deterministic bounds**:
 - (virtual) delay bound
 - backlog bounds
 - departure bound



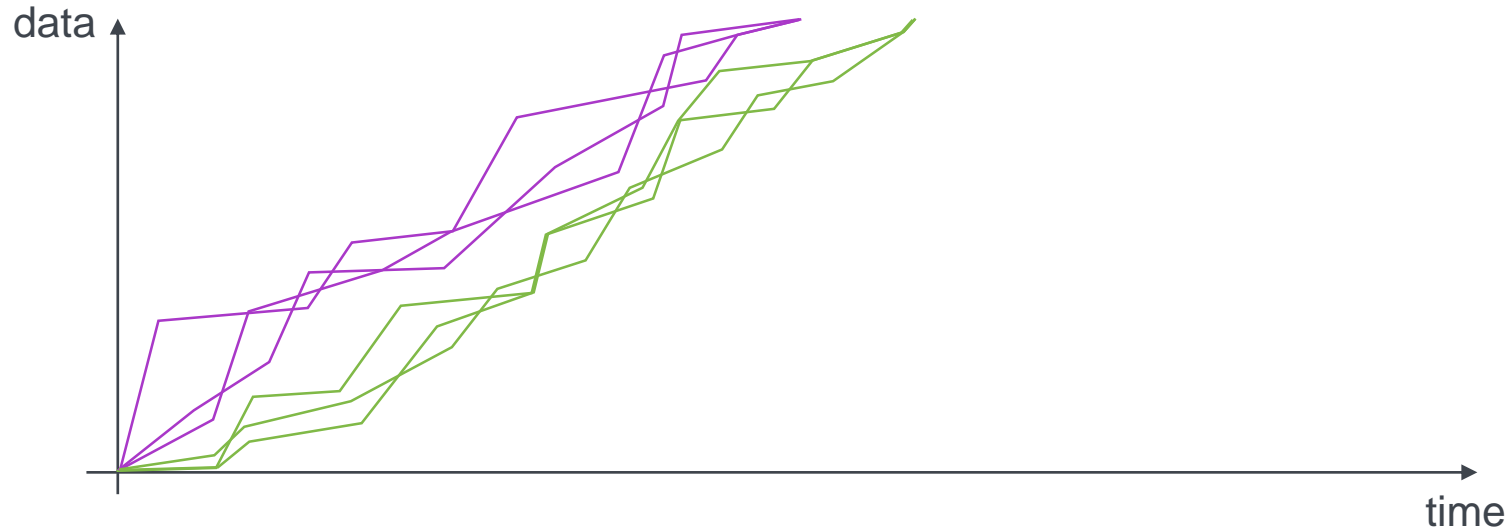
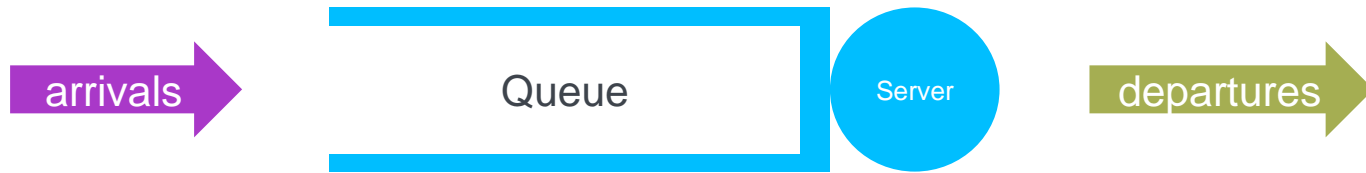
Network calculus

Operating on cumulative curves.



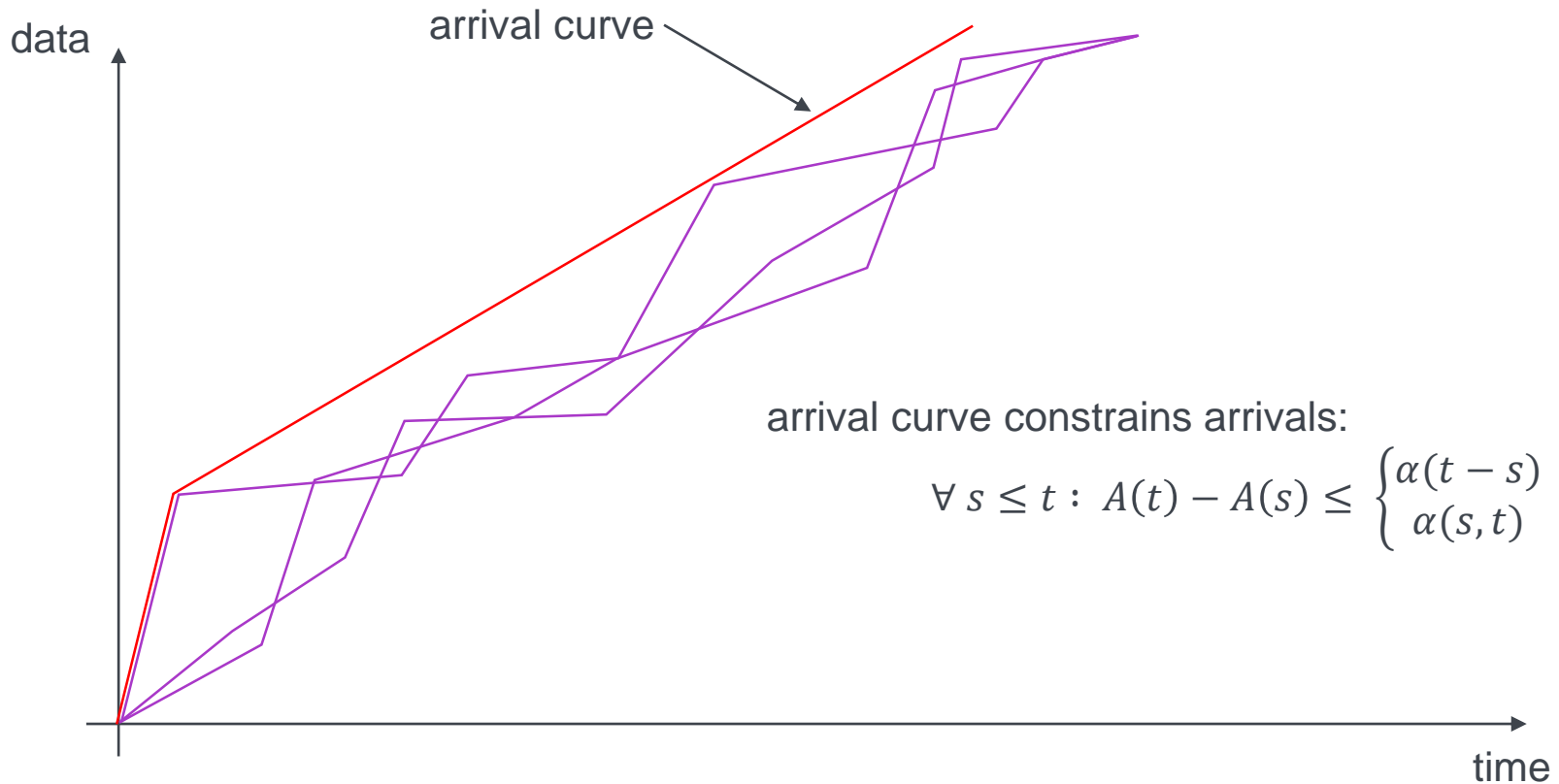
Network calculus

Consider the worst-case (deterministic network calculus).



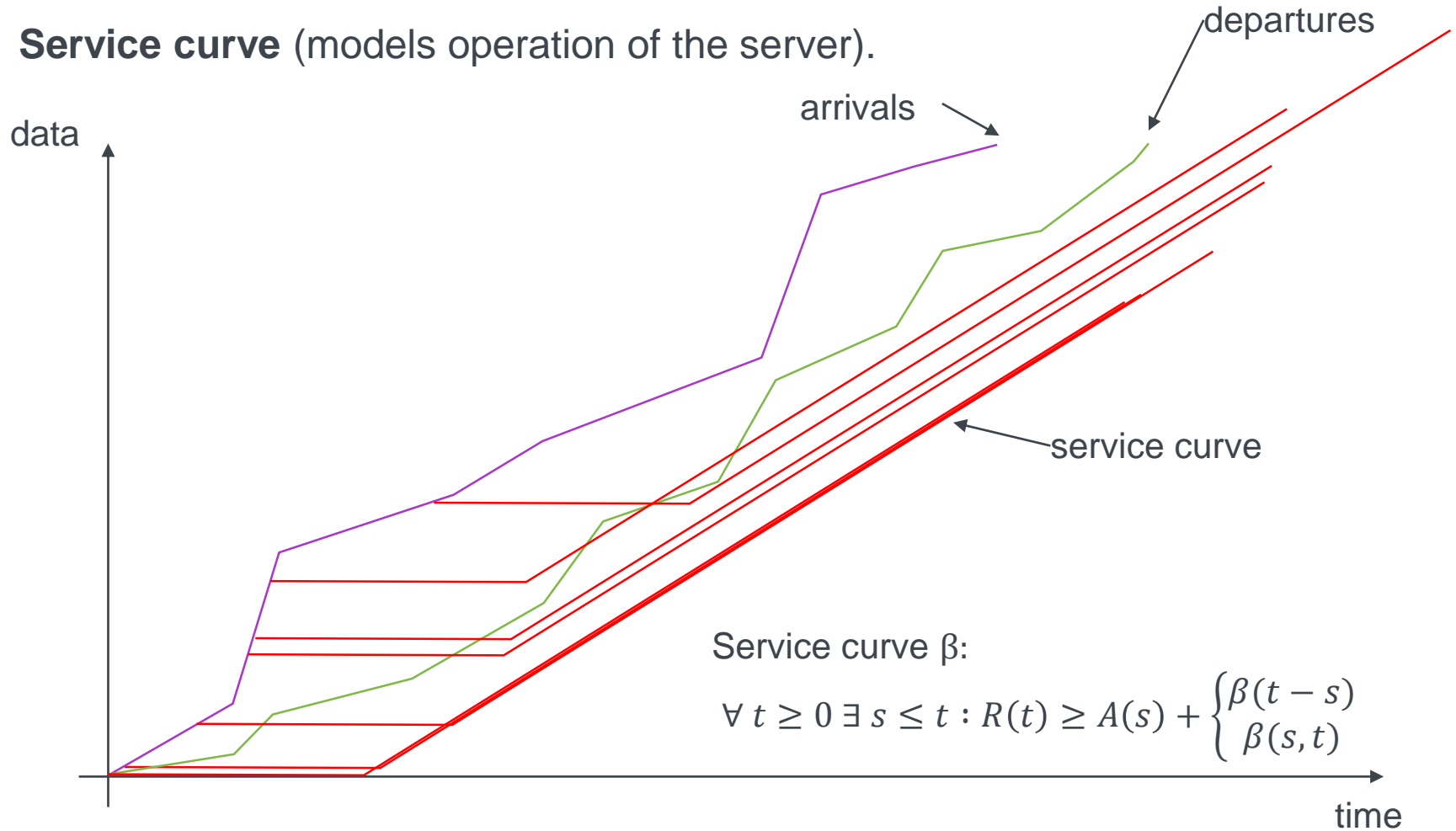
Network calculus

Arrival curve (models data entering the system).

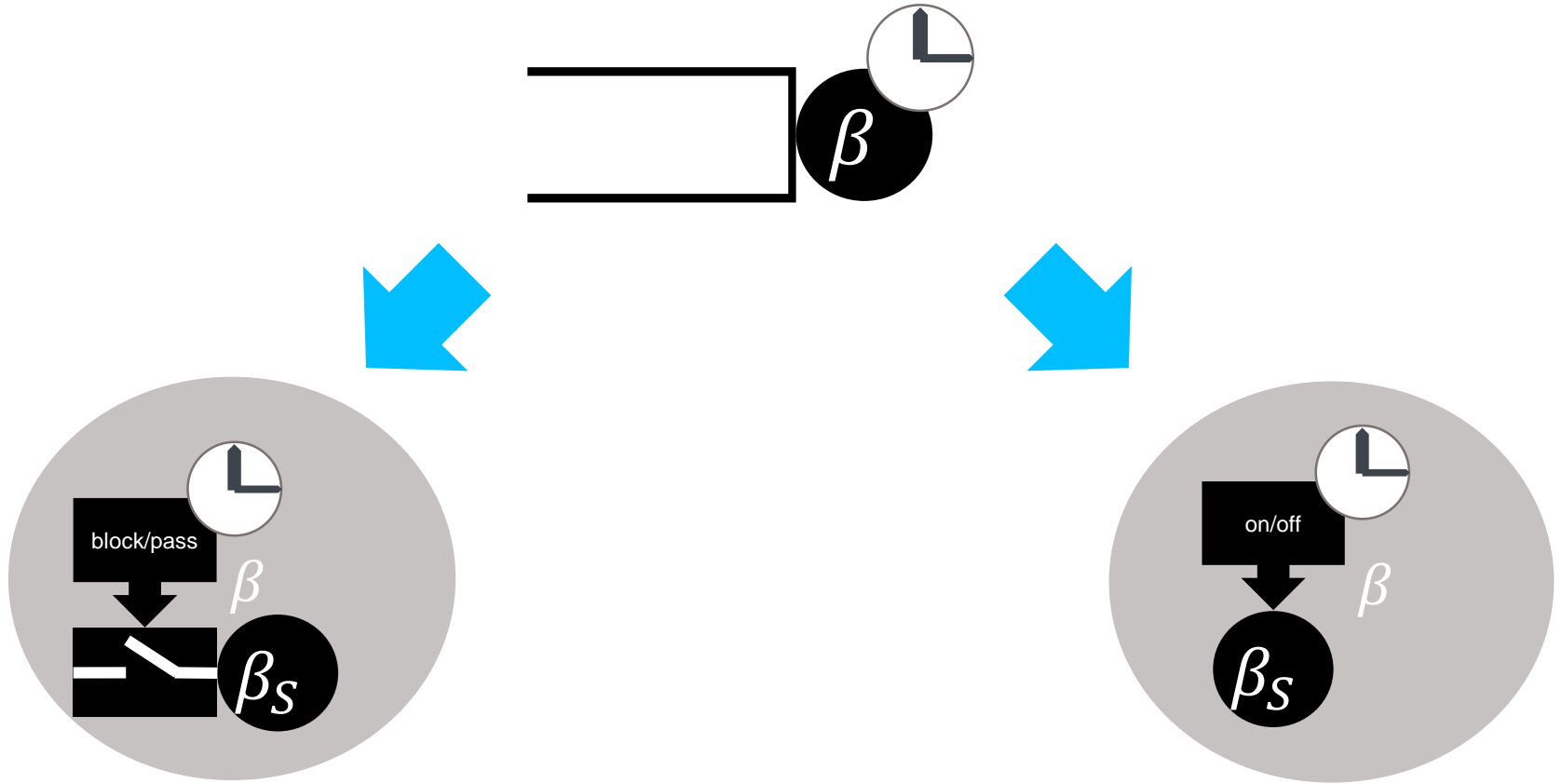


Network calculus

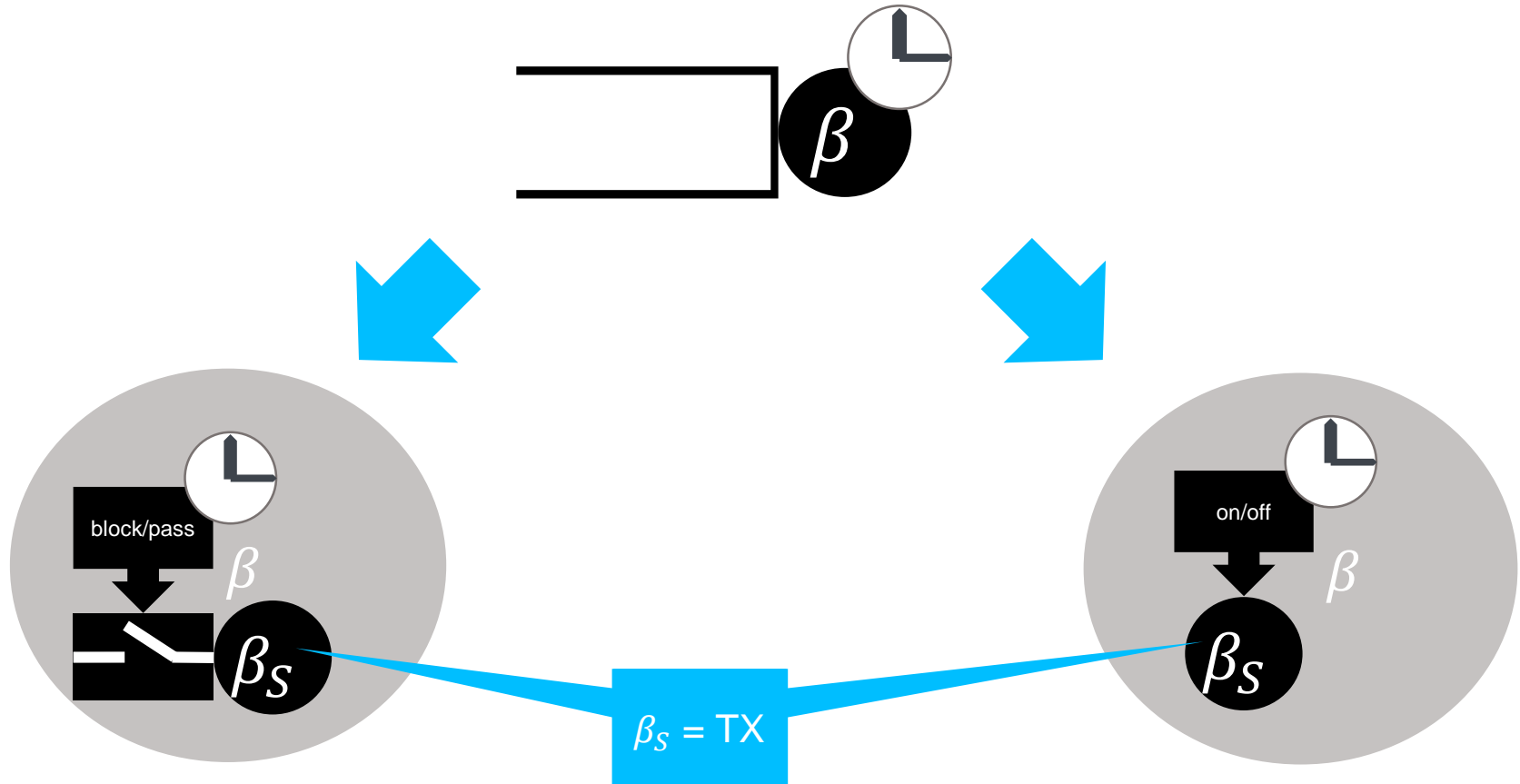
Service curve (models operation of the server).



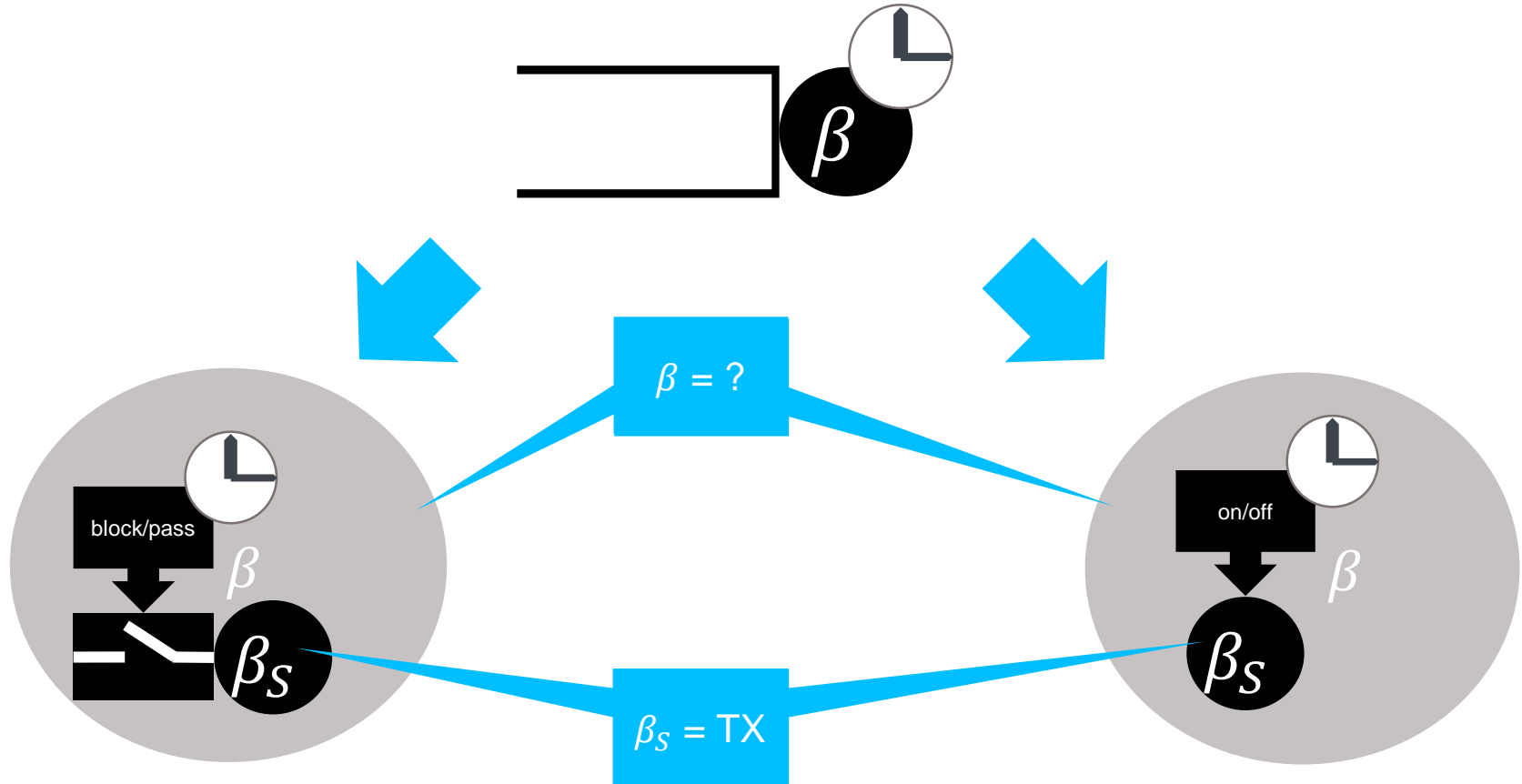
How to model services curves?



How to model services curves?



How to model services curves?



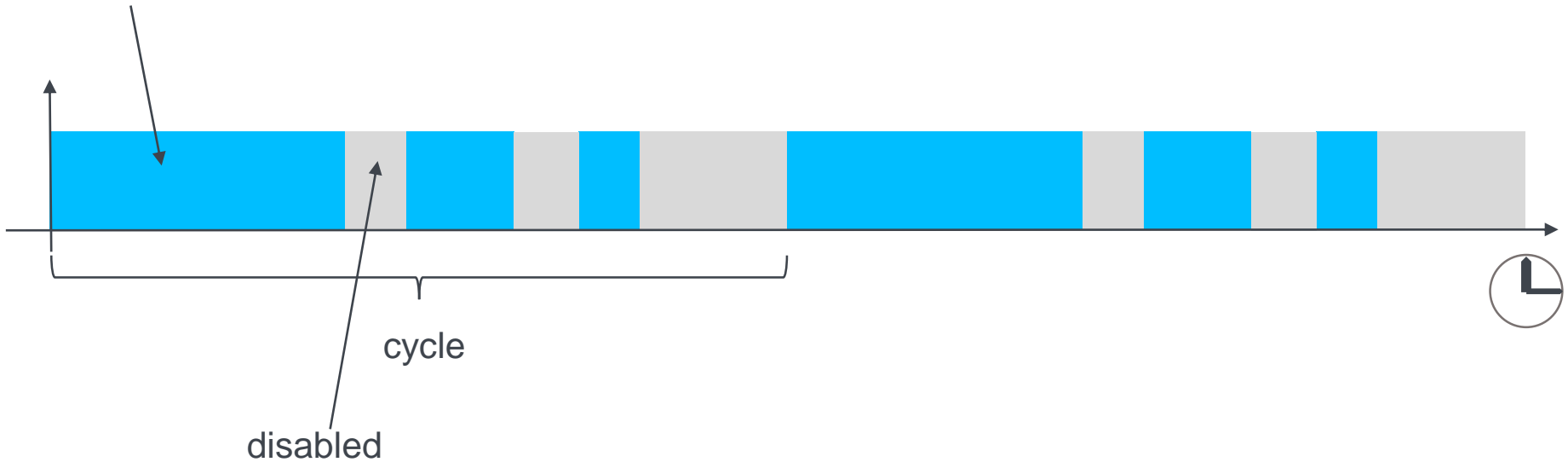
Time-variant service curve

Observation: “enabled” intervals and “disabled” intervals

enabled:

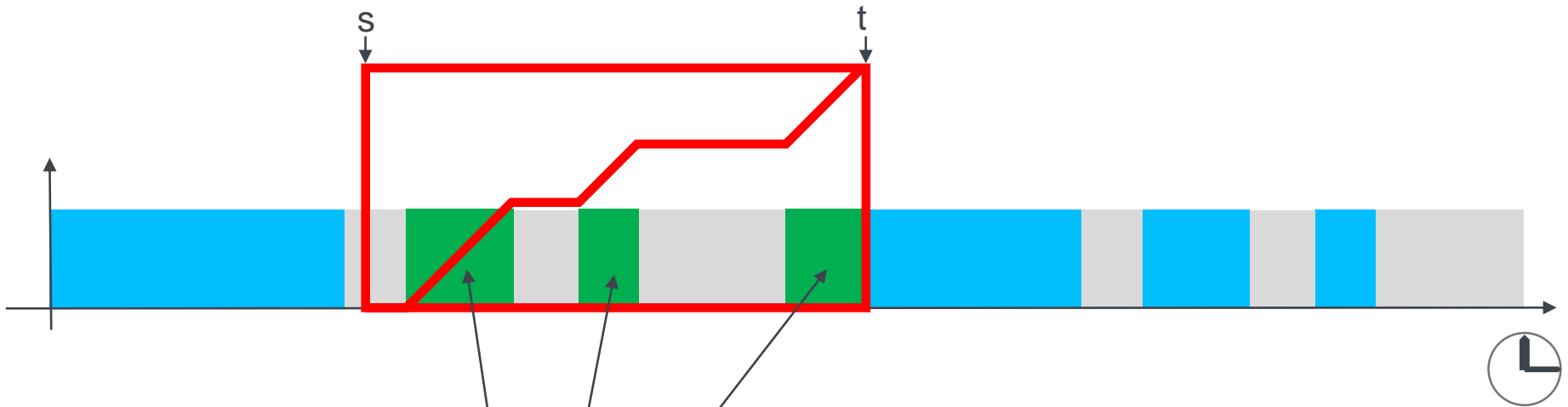
Time-triggered blocking: data is being passed through

Time-triggered halt and restart: server is on



Time-variant service curve

Service in interval $(s,t]$?



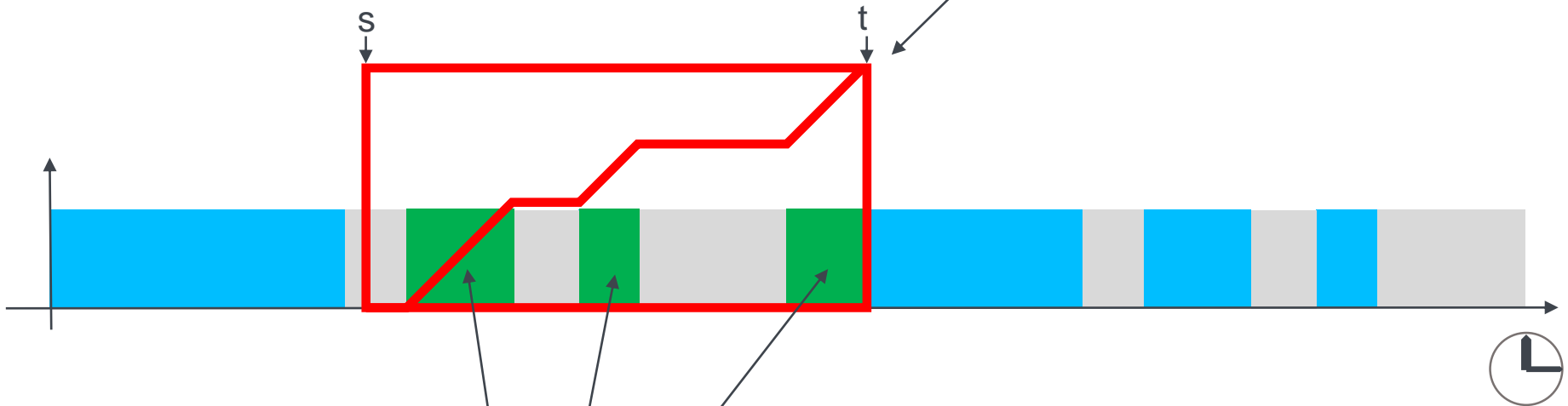
contribute to the total amount of service offered in $(s,t]$

Time-variant service curve

Service in interval $(s,t]$?

Service offered in interval depends on “position” in the schedule

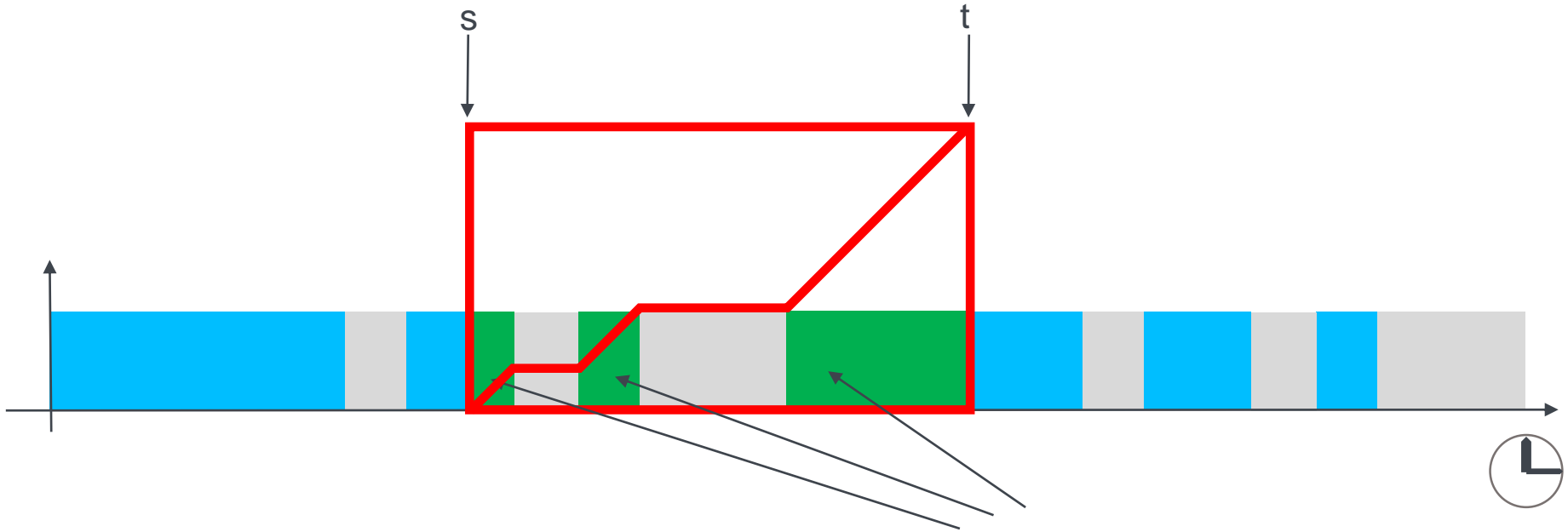
→time-variant



contribute to the total amount of service offered in $(s,t]$

Service curve formulation: idea

Service in interval $(s,t]$?



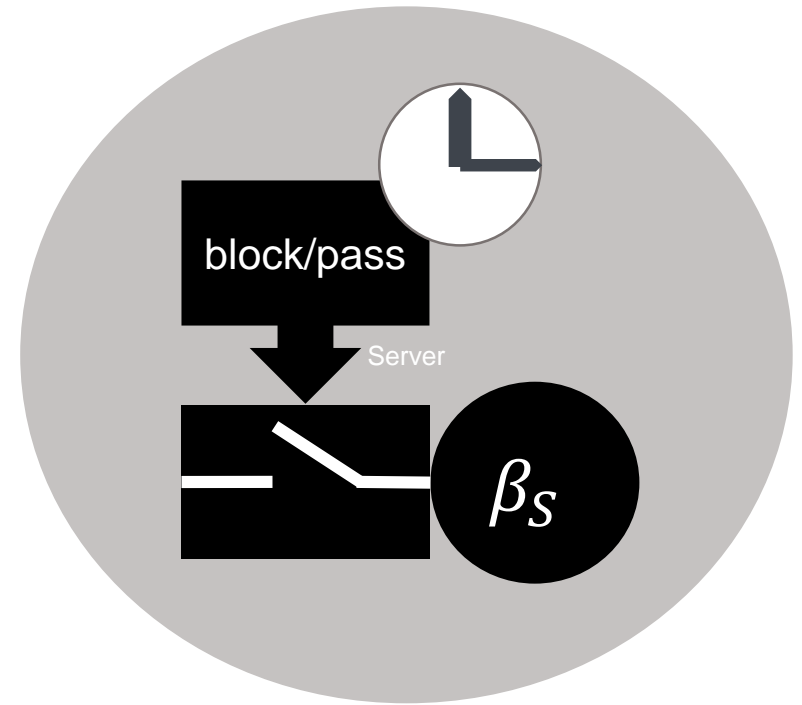
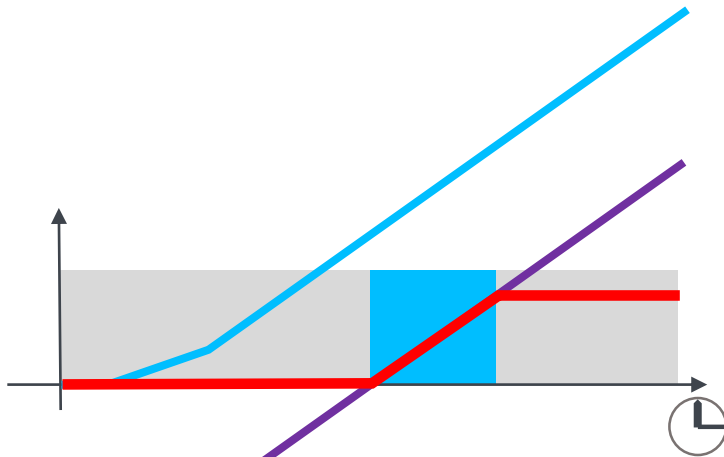
different cases for enabled intervals:

- entirely included / excluded
- partially covered

Time-variant service curve

Time-triggered blocking

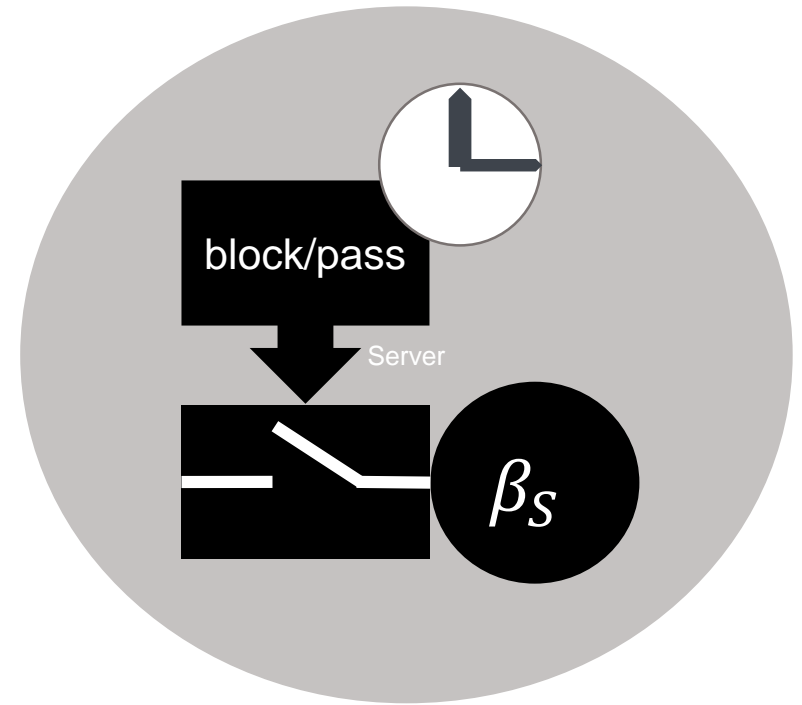
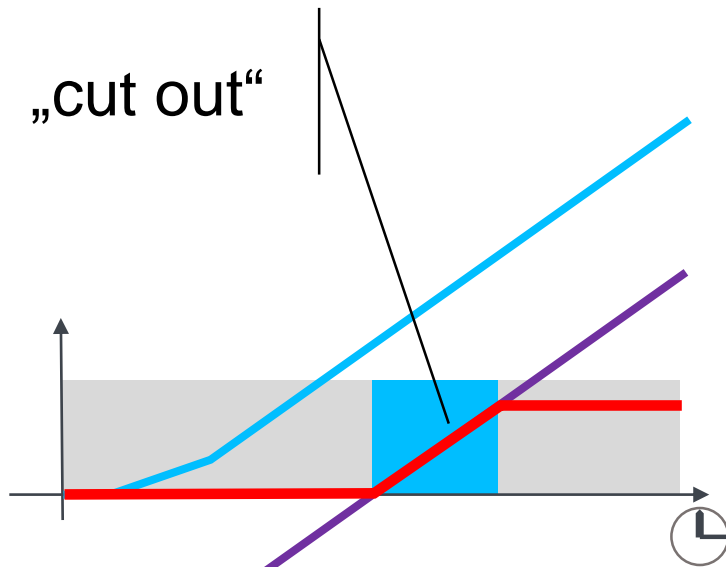
$$\beta(s, t) = \sum_{n=0}^{\infty} [\beta_S(\min(t, t_{d,n})) - \beta_S(\max(s, t_{e,n}))]^+$$



Time-variant service curve

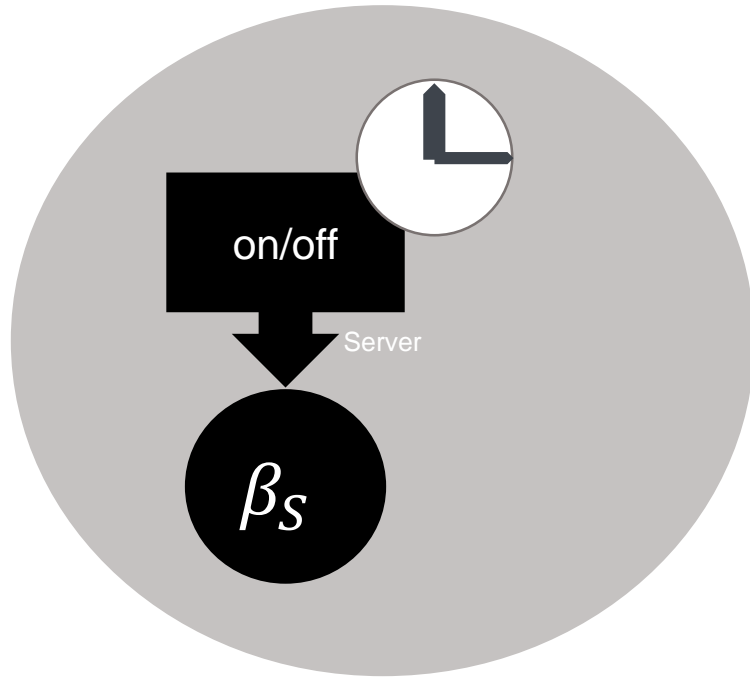
Time-triggered blocking

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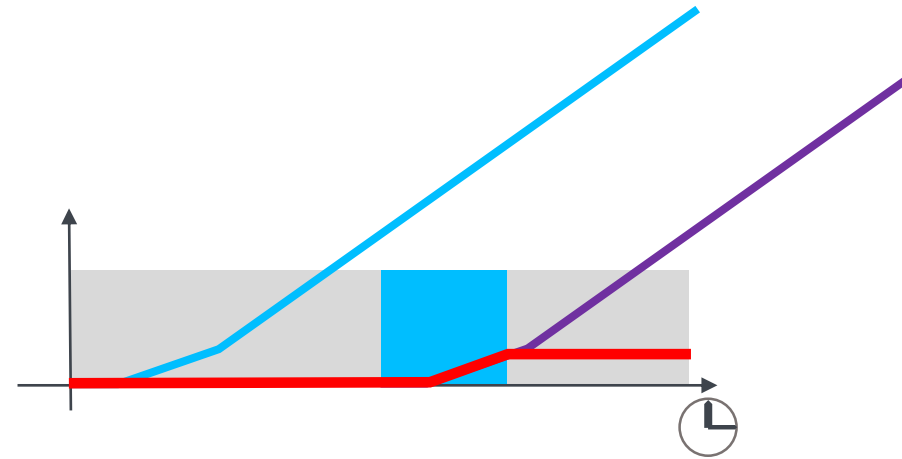


Time-variant service curve

Time-triggered halt and restart

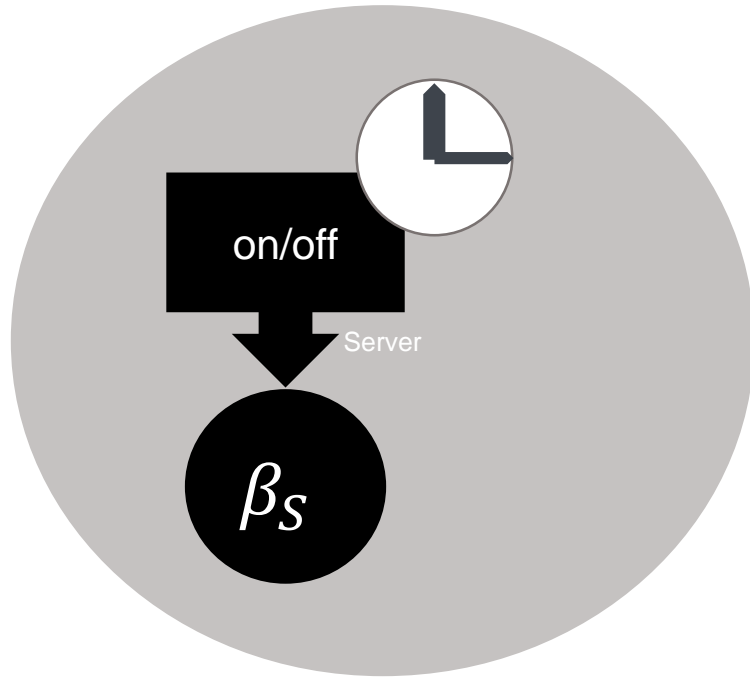


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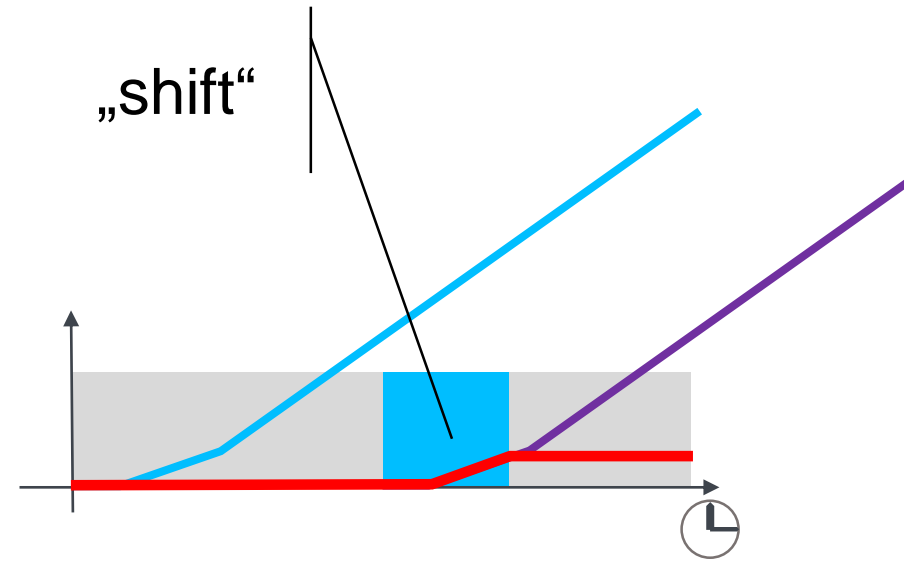


Time-variant service curve

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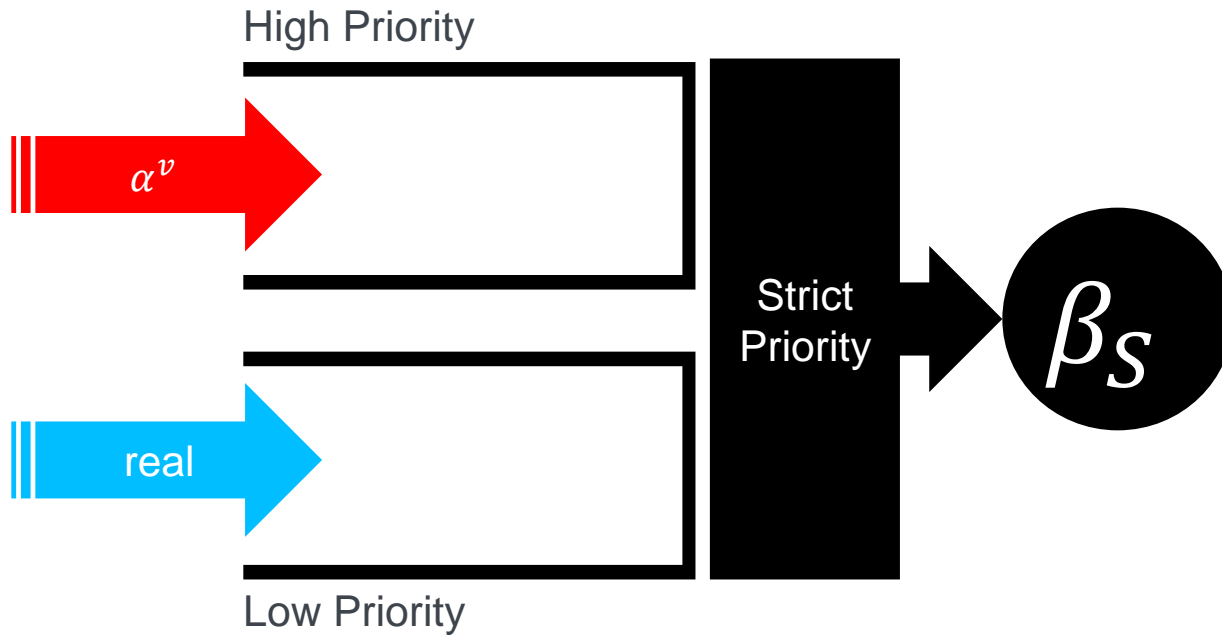


Time-invariant service curve

- Derived from time-variant service curve
- DNC for time-invariant functions
 - “less complicated”
 - computational support available (to some degree)

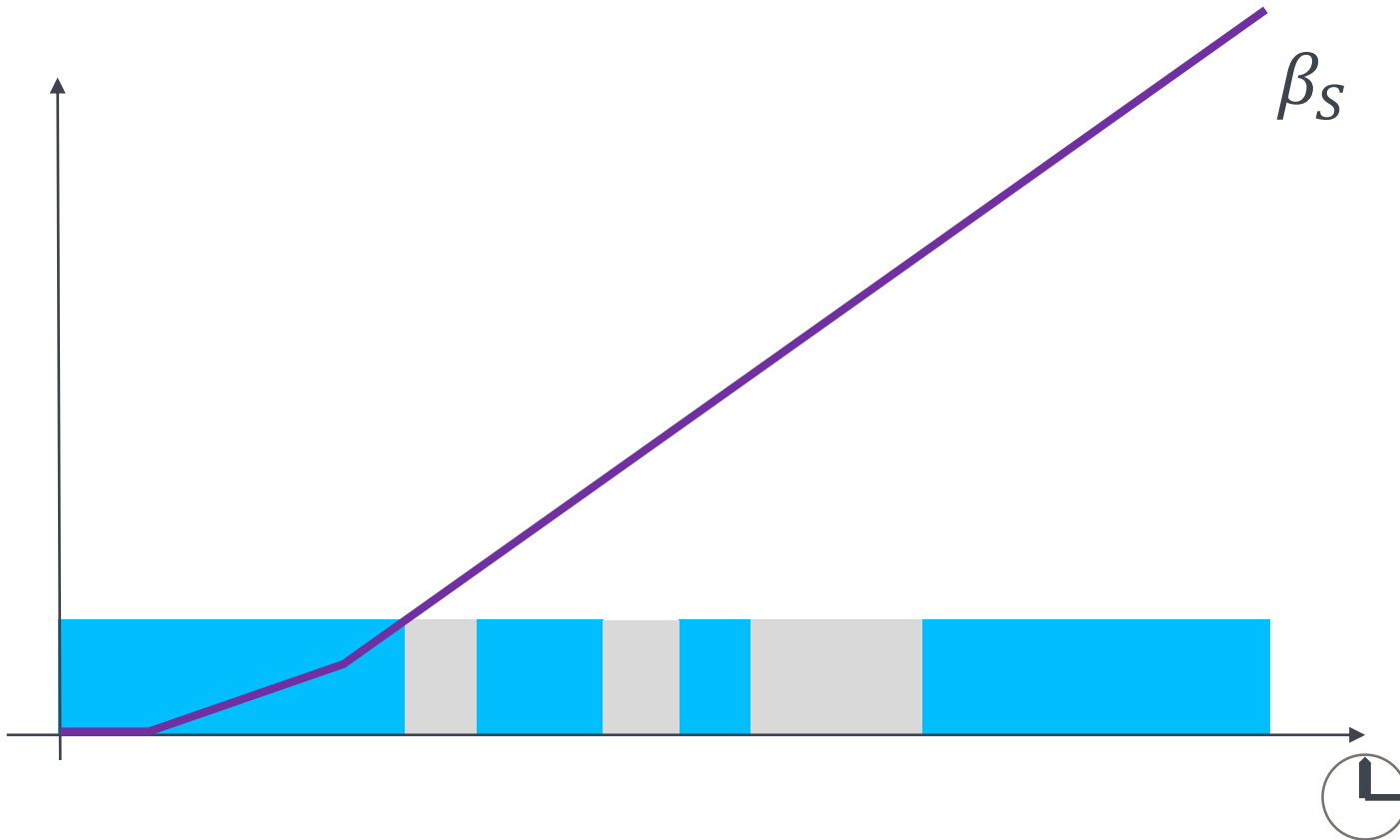
Time-invariant service curve

Leftover service curve: “virtual arrivals”



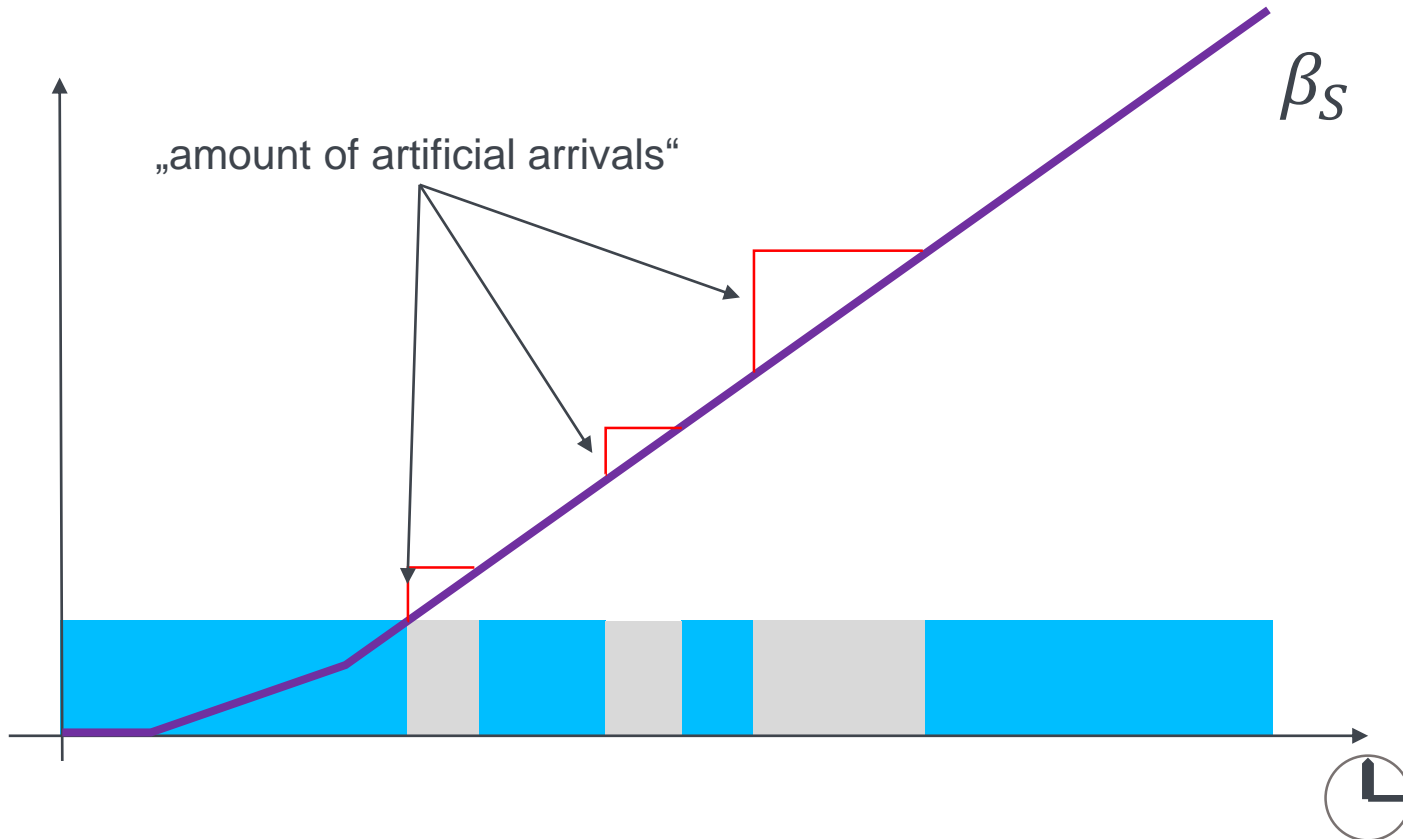
Time-invariant service curve

Leftover service curve: subtract virtual arrivals: $\beta(t) = \left[\sup_{s \leq t} (\beta_S(s) - \alpha^v(s)) \right]^+$



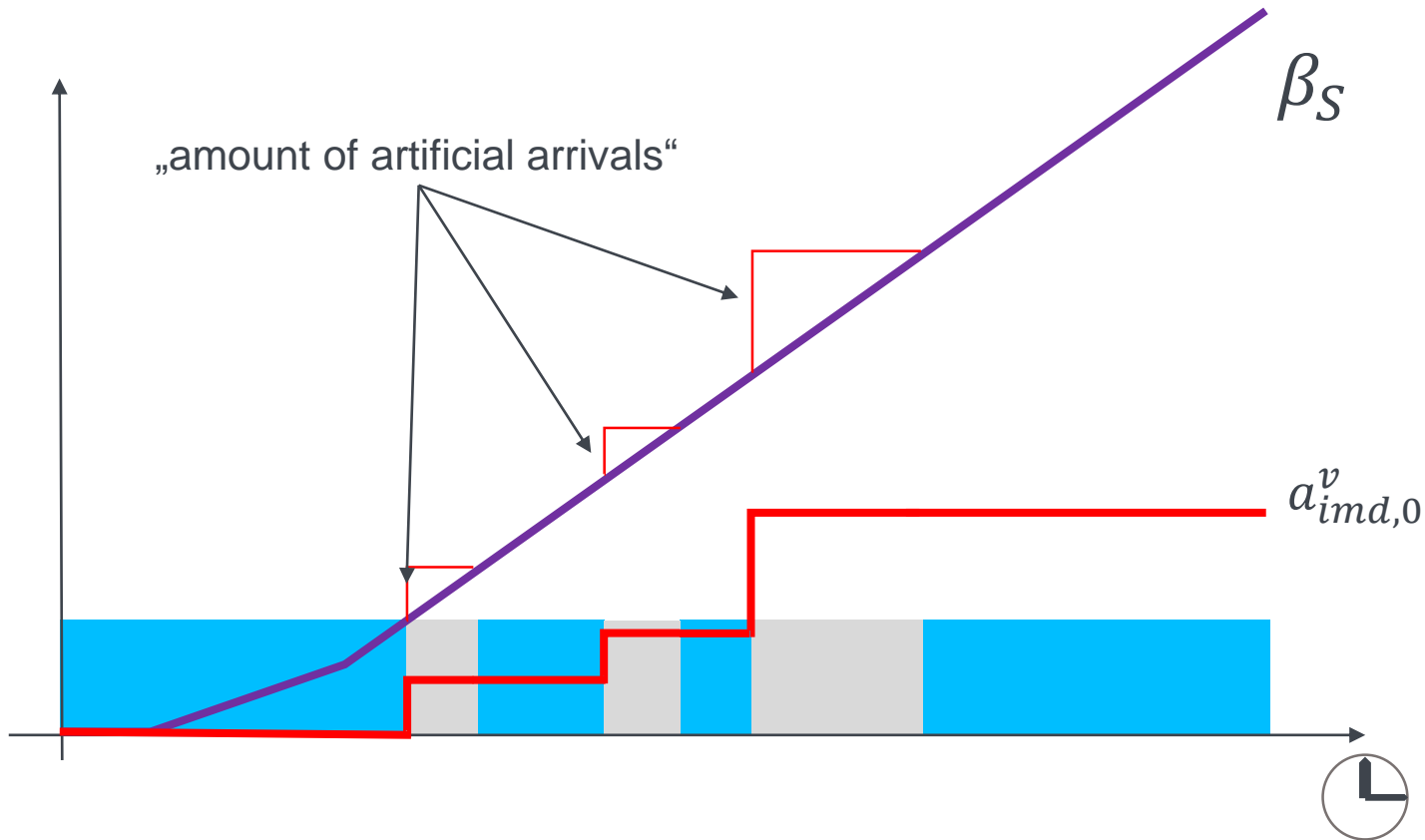
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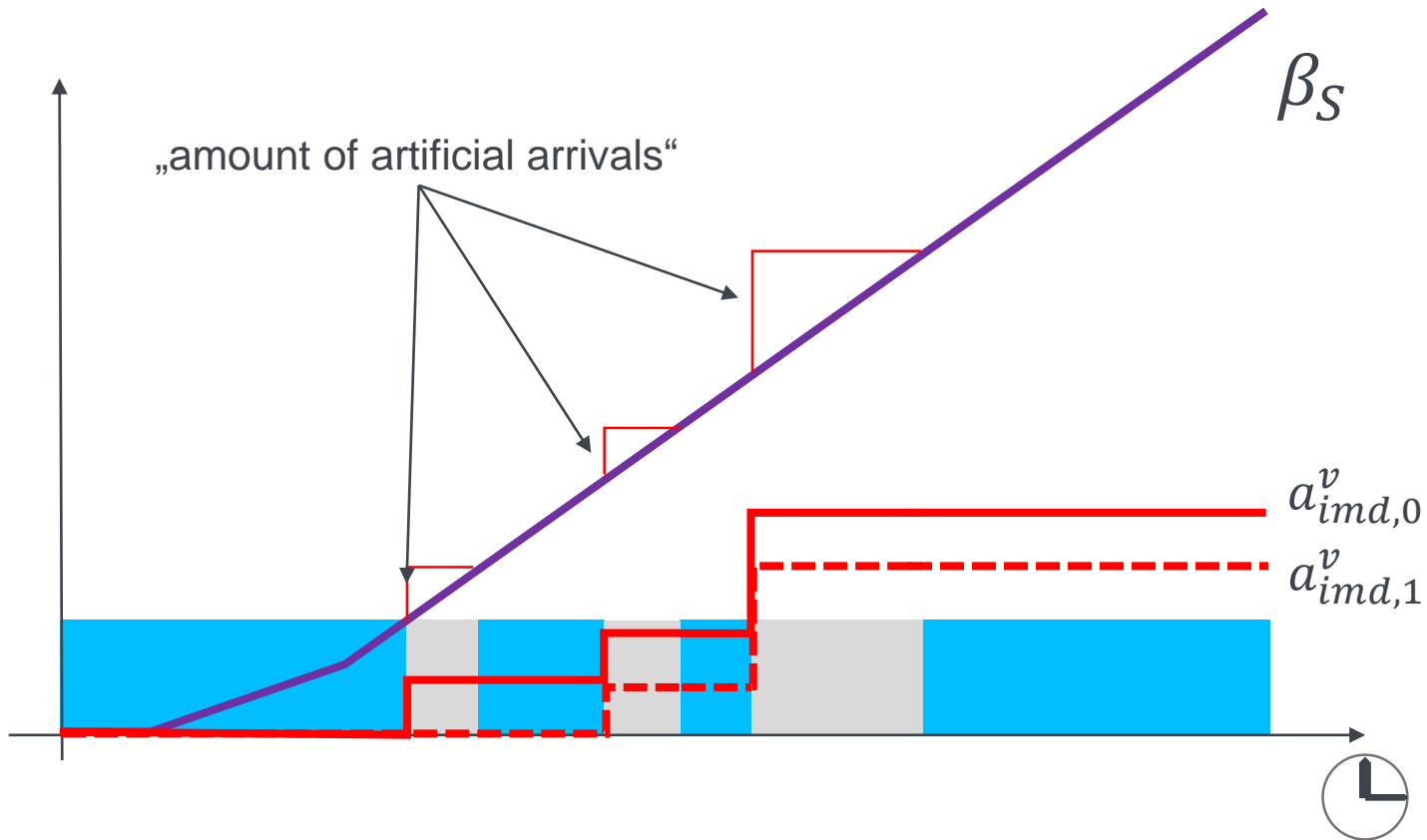
Time-invariant service curve

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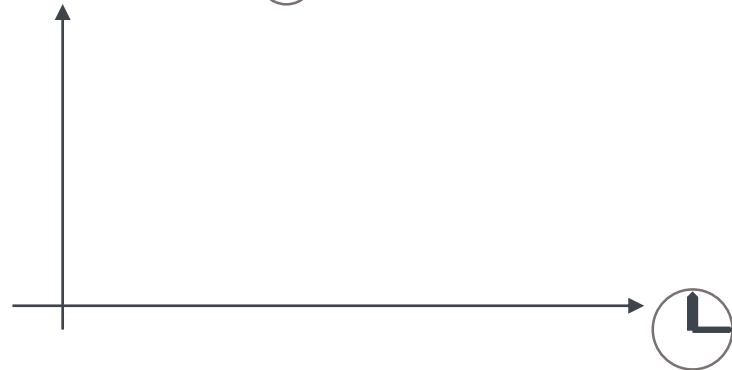
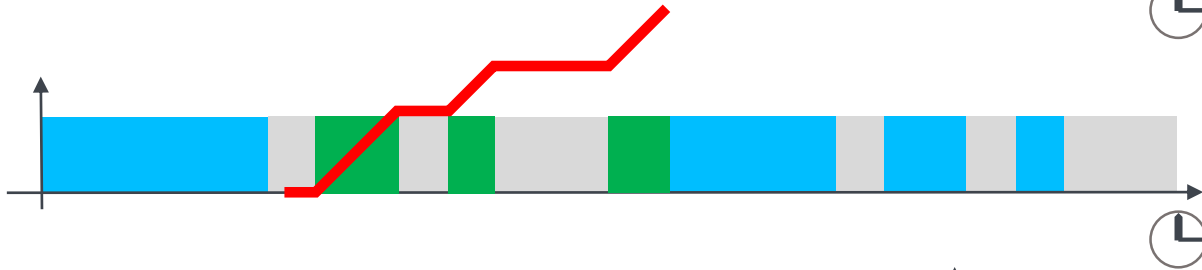
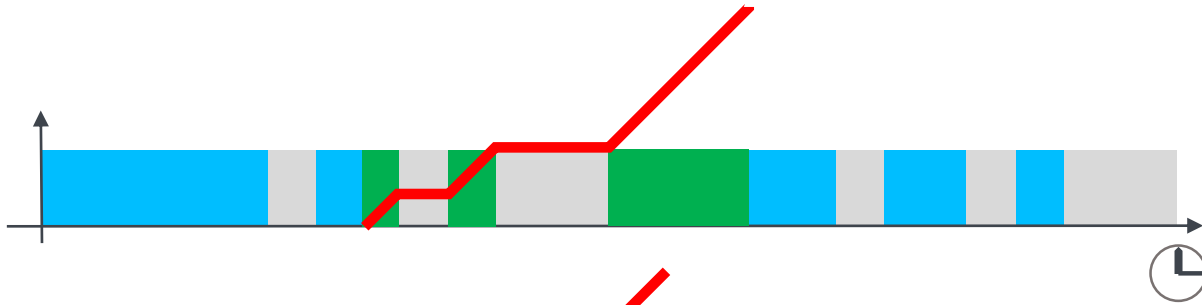
Time-invariant service curve

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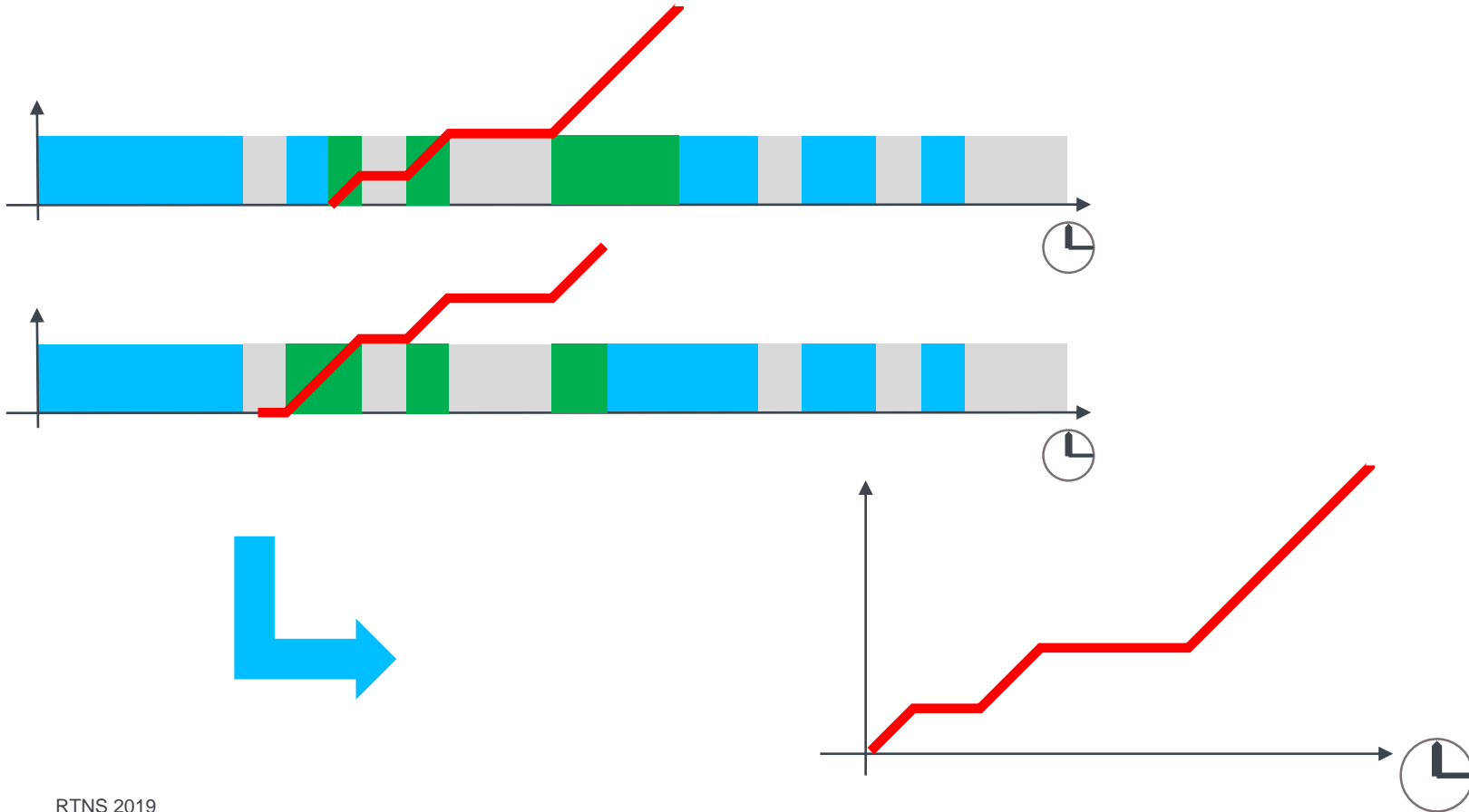
Time-invariant service curve

Direct service curve: find the worst case: $\beta(t) = \inf_{0 < s} (\beta(s, s + t))$



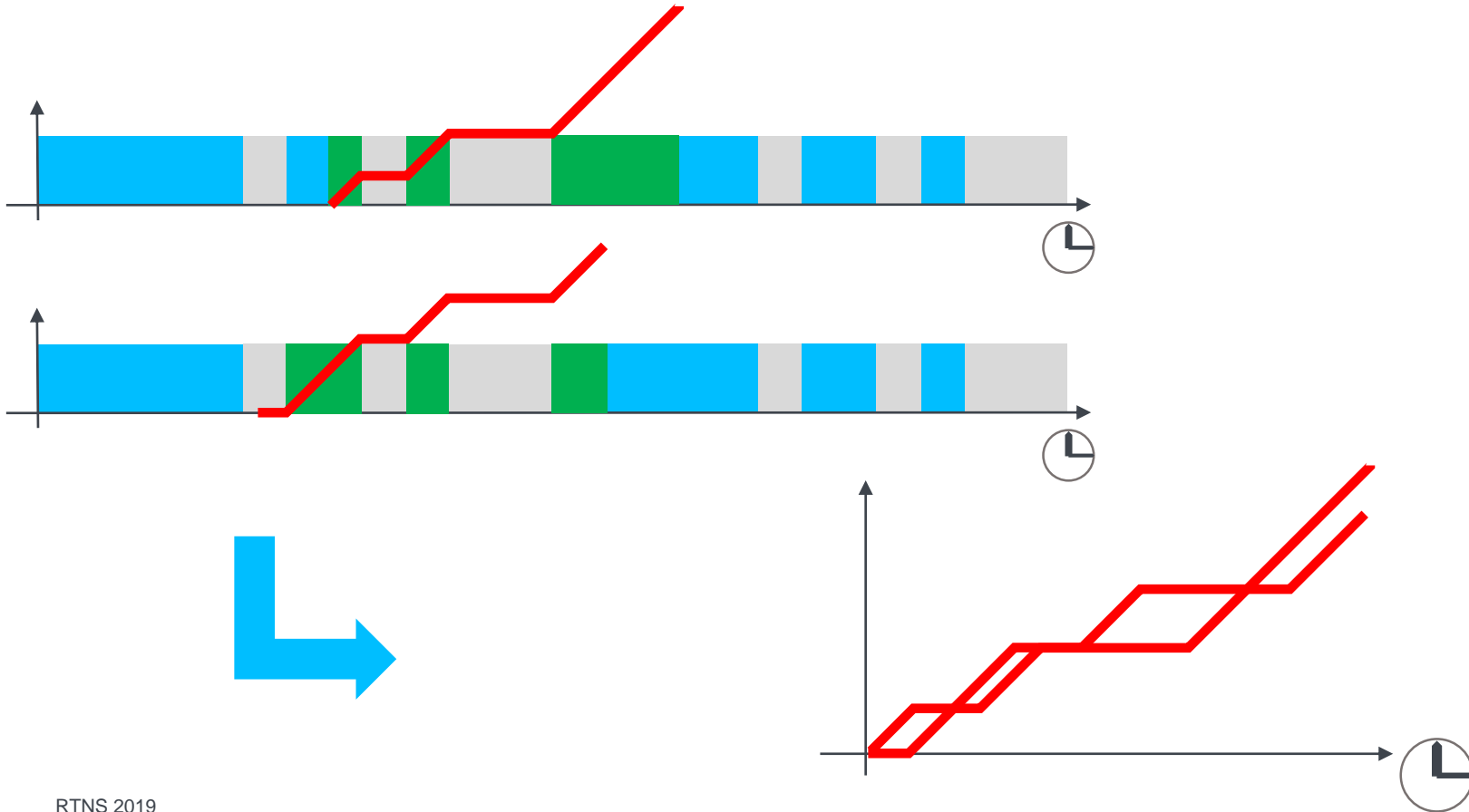
Time-invariant service curve

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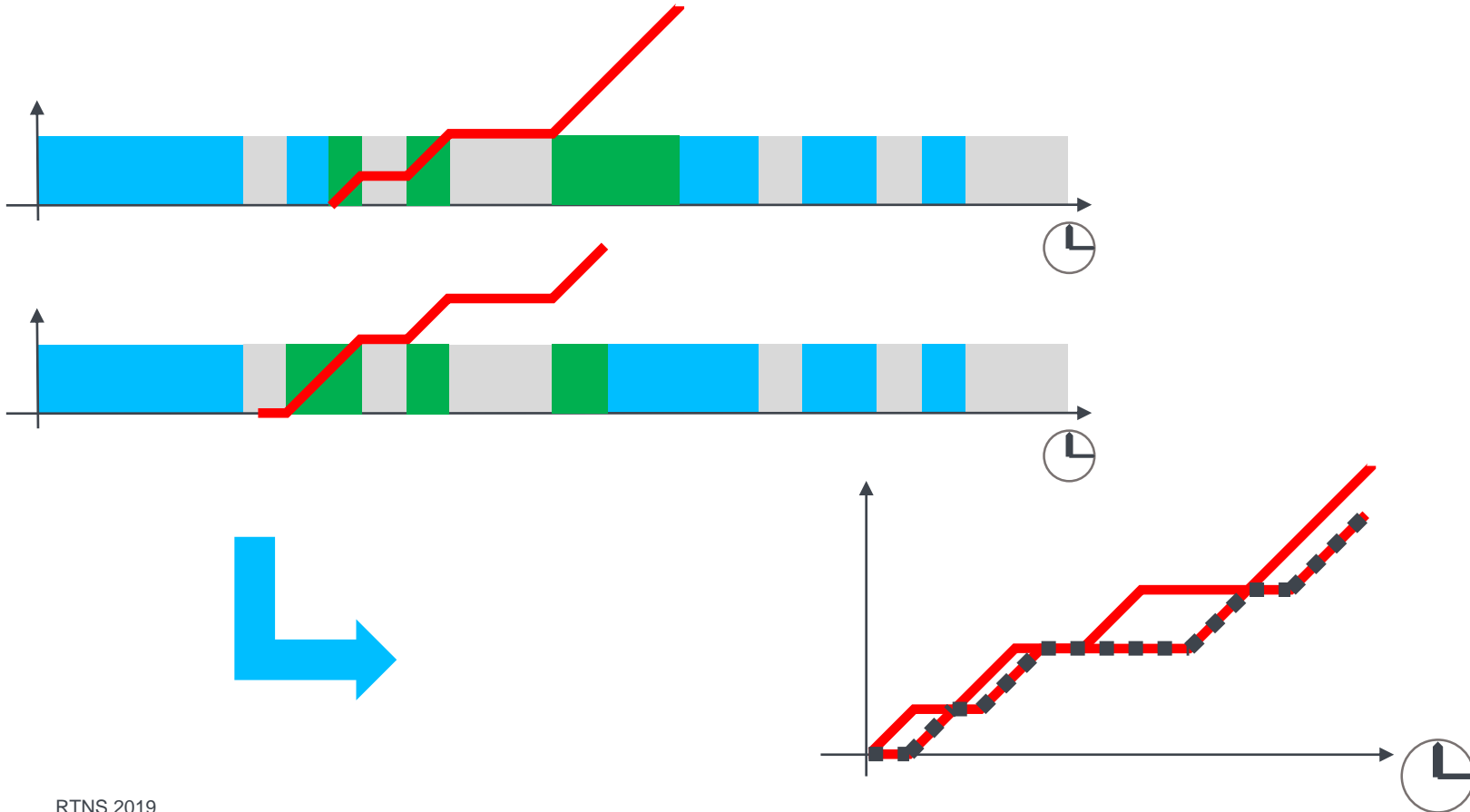
Time-invariant service curve

Direct service curve: find the worst case: $\beta(t) = \inf_{0 < s} (\beta(s, s + t))$



Time-invariant service curve

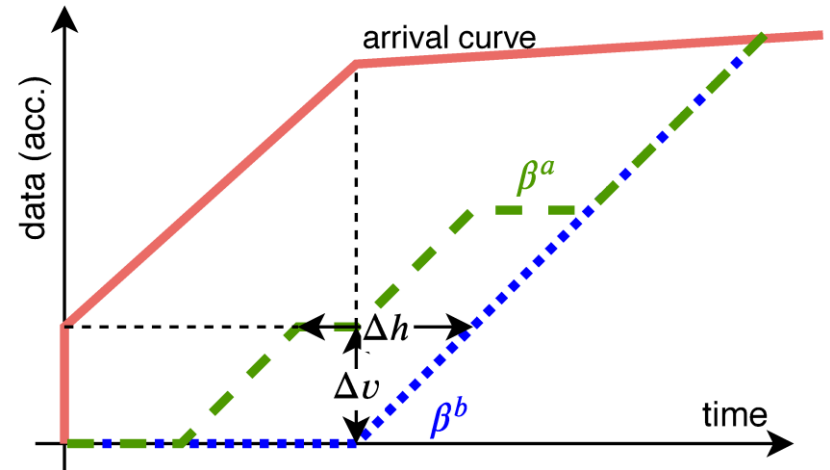
Direct service curve: find the worst case: $\beta(t) = \inf_{0 < s} (\beta(s, s + t))$



Evaluation

How much pessimism?

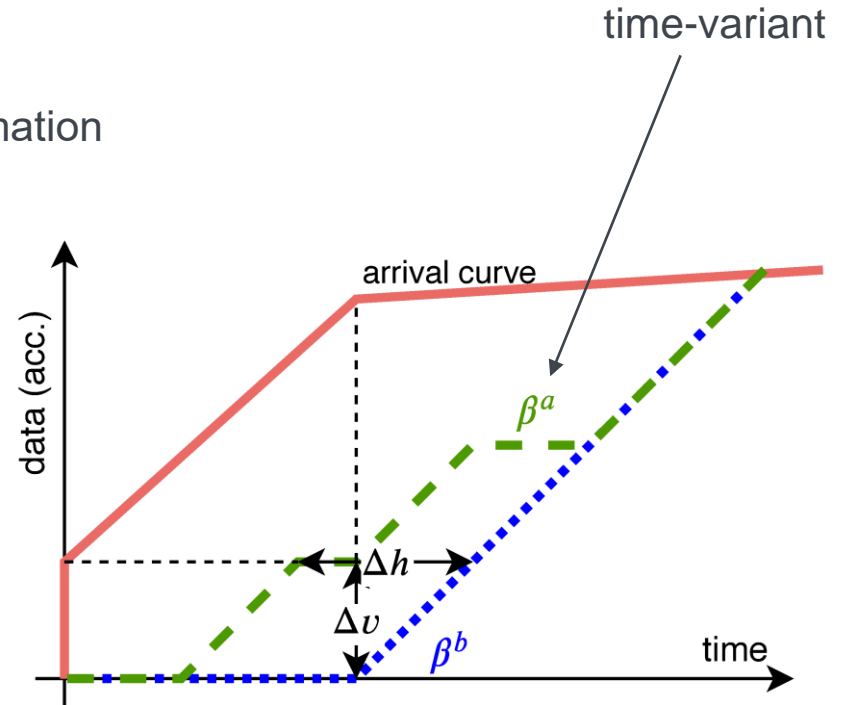
- Randomly generated schedules
- Numerical evaluation of “worst-case” overestimation
- Arrival curve independent metric
 - $\Delta v = \text{backlog}(\beta^a, \beta^b)$
 - $\Delta h = \text{virtual delay}(\beta^a, \beta^b)$



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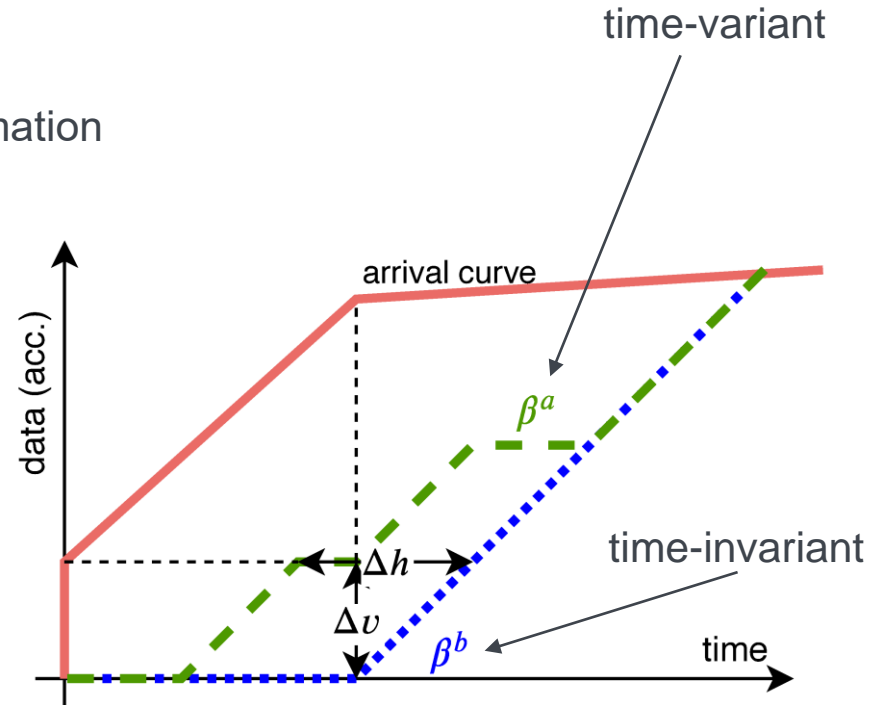
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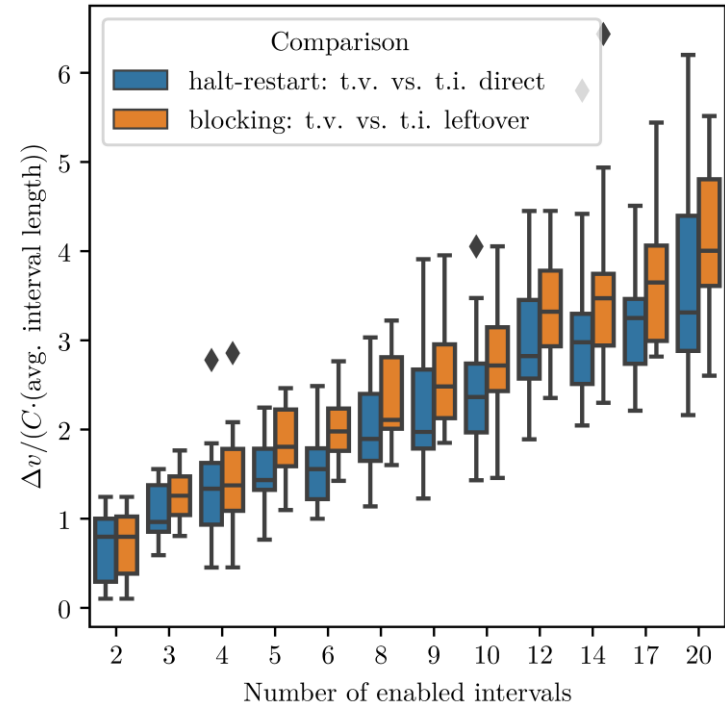
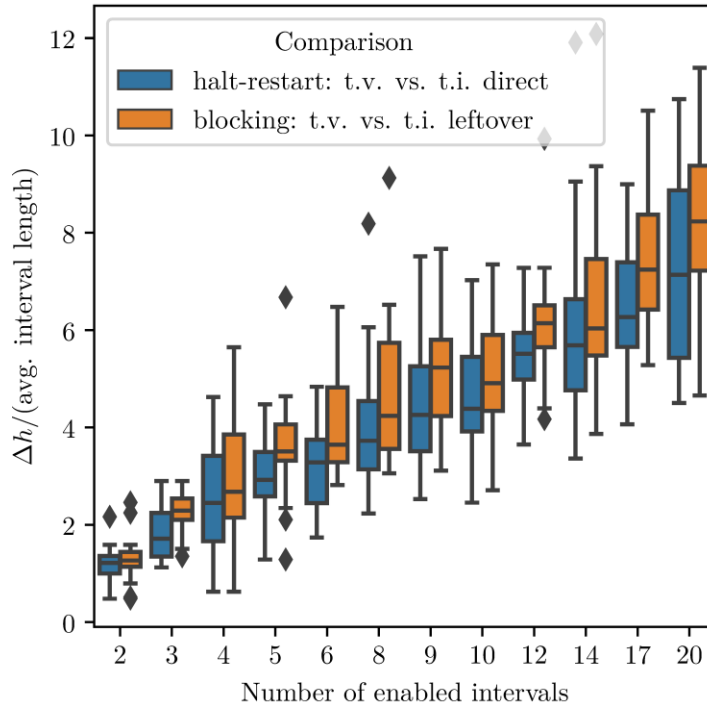


Evaluation

- Proof-of-concept numeric Python implementation
- $\beta_S = t$
- Comparison of time-variant vs. time-invariant for:
 - blocking: t.v. vs. t.i. leftover*
 - halt-restart: t.v. vs. t.i. direct
- Extend t.i. curves with $\beta(s, t) = \beta_{t.i.}(t - s)$

Evaluation results

Random interval length [1,100] time units; 20 schedules per number of enabled intervals

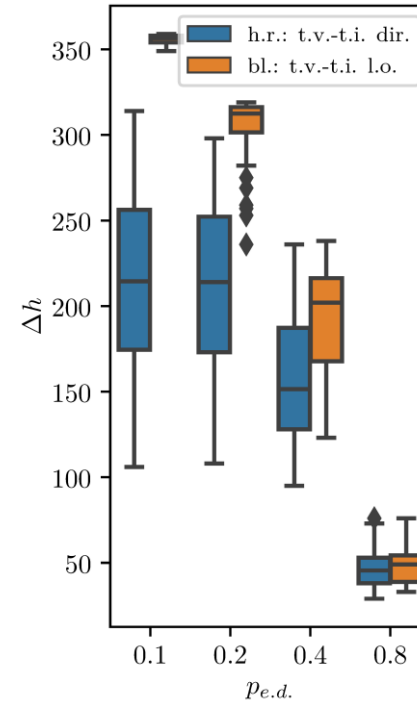


Evaluation results

Varying ratio of enabled/disabled intervals per cycle:
per cycle:

- 10 enabled intervals per cycle
- cycle length: 400 time units

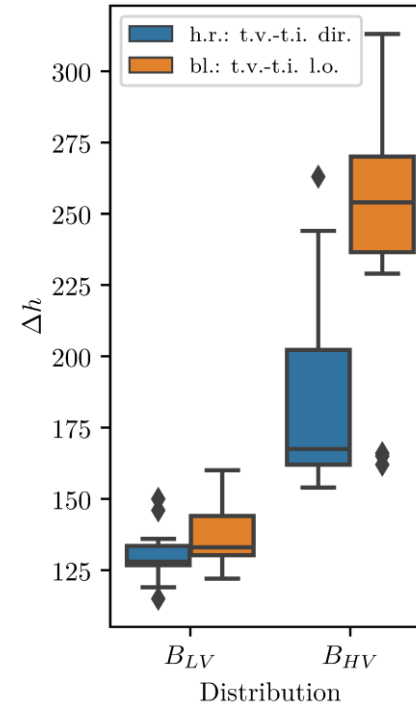
Time-invariant service curve is potentially more pessimistic for long disabled intervals with short interspersed enabled intervals.



Evaluation results

Equal mean interval length, different interval length variance:

- 20 enabled intervals per cycle
- Binomial distribution B_{HV} :
 - Mean interval length: 100
 - Variance: 90
- Binomial distribution B_{LV} :
 - Mean interval length: 100
 - Variance: 20



Discussion

Practical systems have some favorable properties.

- Usually β_S is quite constant
 - Time-invariant service curves can be evaluated more easily
- In converged networks, it is not unlikely to have long enabled intervals with short disabled intervals (from the perspective of the traffic to be analyzed)

Concluding remarks

- Closer look at the fundamental properties of systems with intermittent service
- Similar but Distinct: Blocking vs. Halt-Restart
- Open problems
 - computational support for NC with complicated functions
 - Schedule for service intermittence results in complicated service curves
 - Computational algorithms for time-variant network calculus
 - Multiplexing (i.e., multiple streams sharing one queue)

