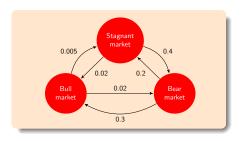
#### James Worrell

Department of Computer Science, Oxford University

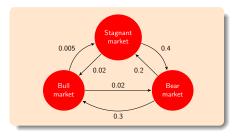
(Joint work with Ventsislav Chonev and Joël Ouaknine)

FICS 2015 September 12th, 2015

# Reachability for Continuous-Time Markov Chains



#### Reachability for Continuous-Time Markov Chains

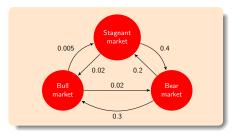


Distribution P(t) at time t satisfies P'(t) = P(t)Q, where

$$Q = \begin{pmatrix} -0.025 & 0.02 & 0.005 \\ 0.3 & -0.5 & 0.2 \\ 0.02 & 0.4 & -0.42 \end{pmatrix}$$

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 Reduce to the time-bounded case by computing the stationary distribution:

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- Reduction does not work if  $\pi$  is on boundary of the target set.
- Semantic shift consider whether the closure of the orbit  $\{P(t): t \geq 0\}$  meets the closure of the target set.

#### Cyber-Physical Systems

"To analyze a cyber-physical system, such as a pacemaker, we need to consider the **discrete software controller** interacting with the physical world, which is typically modelled by **differential equations**"

Rajeev Alur (CACM, 2013)



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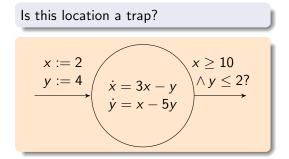
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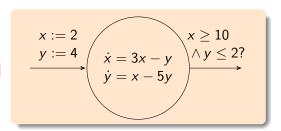
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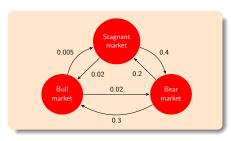
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Is this location a trap?





Is ever more likely to be a Bear market than a Bull market:

$$\exists t (P(t)_{\text{Bear}} \geq P(t)_{\text{Bull}}) ?$$

 $\mathbf{x}: \mathbb{R}_{\geq 0} \to \mathbb{R}^k$ 

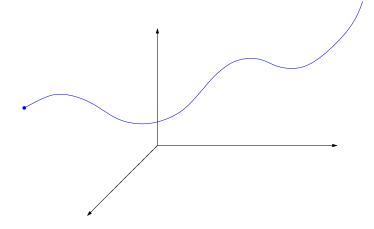
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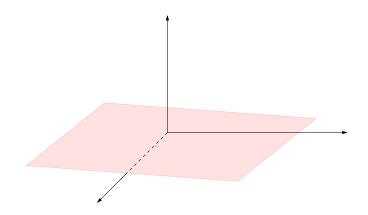
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\Rightarrow \mathbf{x}(t) = \exp(\mathbf{A}t)\mathbf{x}(0)
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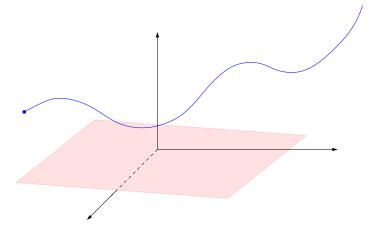
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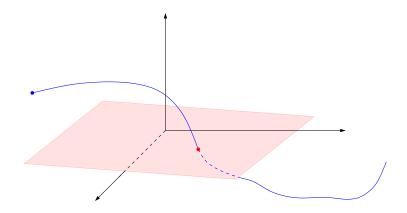
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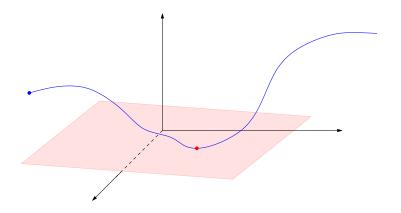
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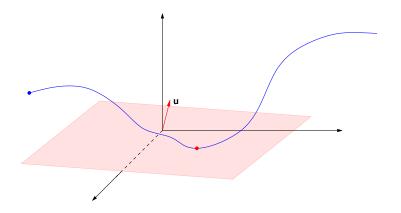
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Note – the  $\lambda_j$  are complex in general.

Let  $f: \mathbb{R}_{\geq 0} \to \mathbb{R}$  be given as above, with all coefficients algebraic.

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#### **BOUNDED-ZERO Problem**

Instance: f and bounded interval [a, b]

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Question: Is there  $t \in \mathbb{R}_{\geq 0}$  such that f(t) = 0?

• Decidability open! [Bell, Delvenne, Jungers, Blondel 2010]

#### Related Work

A lot of work since 1920s on the zeros of exponential polynomials

$$f(z) = \sum_{j=1}^{m} P_j(z) e^{\lambda_j z}$$

(Polya, Ritt, Tamarkin, Kac, Voorhoeve, van der Poorten, ...) but mostly on distribution of *complex* zeros.

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#### **CONTINUOUS-ORBIT Problem**

The problem of whether the trajectory  $\mathbf{x}(t) = e^{\mathbf{A}t}\mathbf{x}(0)$  reaches a given target point was shown to be decidable by Hainry (2008) and in PTIME by Chen, Han and Yu (2015).

Theorem (Bell, Delvenne, Jungers, Blondel 2010)

In dimension 2, BOUNDED-ZERO and ZERO are decidable.

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Assuming Schanuel's Conjecture, BOUNDED-ZERO is decidable in all dimensions.

It turns out that this result (in fact, a powerful generalisation of it) had already been discovered (but never published) in the early 1990s by Macintyre and Wilkie!

[Angus Macintyre, personal communication, July 2015]

Theorem (Chonev, Ouaknine, W. 2015)

In dimension 8 or less, ZERO reduces to BOUNDED-ZERO.

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In dimension 9 (and above), decidability of ZERO would entail major breakthroughs in Diophantine approximation—the Diophantine approximation type of  $\alpha$  would be computable to within arbitrary precision.

## Schanuel's Conjecture

### Theorem (Lindemann-Weierstrass)

If  $a_1, \ldots, a_n$  are algebraic numbers linearly independent over  $\mathbb{Q}$ , then  $e^{a_1}, \ldots, e^{a_n}$  are algebraically independent.

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If  $z_1,\ldots,z_n$  are complex numbers linearly independent over  $\mathbb Q$  then some n-element subset of  $\{z_1,\ldots,z_n,e^{z_1},\ldots,e^{z_n}\}$  is algebraically independent.

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#### Example

By Schanuel's conjecture some two-element subset of  $\{1,\pi i,e^1,e^{\pi i}\}$  is algebraically independent.

Real-valued exponential polynomial 
$$f(t) = \sum_{j=1}^m P_j(t) e^{\lambda_j t}$$

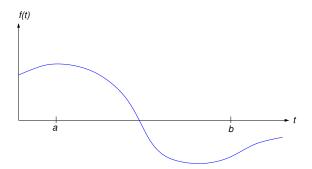
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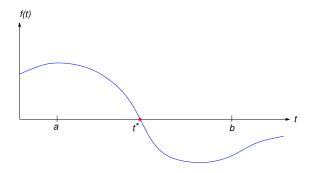
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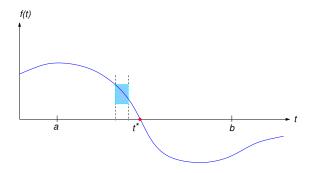


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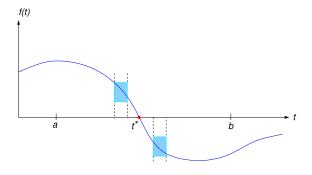
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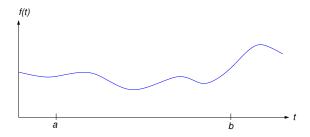
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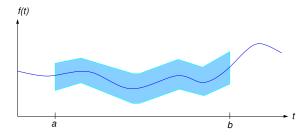


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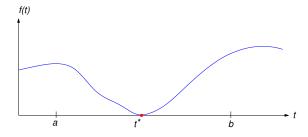
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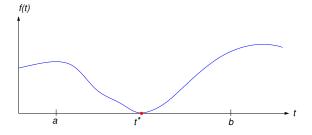
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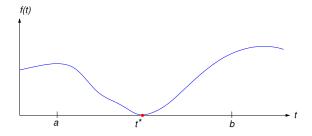


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Can this situation arise?

Real-valued exponential polynomial  $f(t) = \sum_{j=1}^{m} P_j(t) e^{\lambda_j t}$ 



Easily! For example,  $f(t) = 2 + e^{it} + e^{-it}$ .

#### Example

• Write  $f(t) = 2 + e^{it} + e^{-it}$  in the form  $f(t) = P(e^{it})$  for the Laurent polynomial

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**Idea**: factorise f. Noting that factors may be complex-valued!

### The Real Case

Any exponential polynomial f(t) can be written

$$f(t) = P(t, e^{a_1 t}, \dots, e^{a_m t})$$

with

$$P \in \mathbb{C}[x, x_1^{\pm 1}, \dots, x_m^{\pm 1}]$$

and  $\{a_1, \ldots, a_m\}$  a set of real and imaginary algebraic numbers that is linearly independent over  $\mathbb{Q}$ .

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Complex case requires some new ideas . . .

### The Unbounded Case

"there are known unknowns; that is to say we know there are some things we do not know."



### Continued Fractions

Finite continued fractions:

$$[3,7,15,1,292] = 3 + \frac{1}{7 + \frac{1}{15 + \frac{1}{1 + \frac{1}{297}}}}$$

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$$= 3.141592653...$$

Infinite continued fractions:

$$[a_0, a_1, a_2, a_3, \ldots] = a_0 + \frac{1}{a_1 + \frac{1}{a_2 + \frac{1}{a_3 + \cdots}}}$$

#### Theorem

The continued fraction expansion of a real quadratic irrational number is periodic.

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$$\sqrt{389} = [19, 1, 2, 1, 1, 1, 1, 2, 1, 38, 1, 2, 1, 1, 1, 1, 2, 1, 38, \dots]$$

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What about numbers of degree  $\geq$  3?

$$\sqrt[3]{2} = [1, 3, 1, 5, 1, 1, 4, 1, 1, 8, 1, 14, 1, 10, 2, 1, 4, 12, 2, 3, 2, 1 \\ 3, 4, 1, 1, 2, 14, 3, 12, 1, 15, 3, 1, 4, 534, 1, 1, 5, 1, 1, \ldots]$$

# Real Algebraic Numbers

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$$\sqrt{389} = [19, 1, 2, 1, 1, 1, 1, 2, 1, 38, 1, 2, 1, 1, 1, 1, 2, 1, 38, \ldots]$$

What about numbers of degree  $\geq$  3?

$$\sqrt[3]{2} = [1, 3, 1, 5, 1, 1, 4, 1, 1, 8, 1, 14, 1, 10, 2, 1, 4, 12, 2, 3, 2, 1 \\ 3, 4, 1, 1, 2, 14, 3, 12, 1, 15, 3, 1, 4, 534, 1, 1, 5, 1, 1, \ldots]$$

Lang and Trotter: "no significant departure from random behaviour"

# An Open Problem

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"Is there an algebraic number of degree higher than two whose simple continued fraction has unbounded partial quotients? Does every such number have unbounded partial quotients?"

R. K. Guy, 2004



# A Mathematical Obstacle at Dimension 9

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#### Remark

Perhaps this set is recursive—it may even be  $\emptyset$  or  $\mathbb{R} \cap \mathbb{A}$ . However proving recursive enumerability would be a significant achievement.

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$$\left|x-\frac{m}{n}\right|$$

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• Relate this to the existence of zeros of order-9 exponential polynomial f(t) with terms  $e^{ixt}$  and  $e^{it}$ .

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Instance: f

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#### **Theorem**

In dimension 8 or less, ZERO reduces to BOUNDED-ZERO.

- In the limit f is either never zero or infinitely often zero, and we decide which is the case.
- Diophantine-approximation bounds play a key role in the proof.

Consider the exponential polynomial

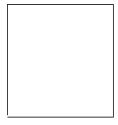
$$f(t) = 1.9 + \cos(t + \varphi_1) + \cos(\sqrt{2}t + \varphi_2) - e^{-t}$$

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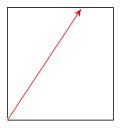
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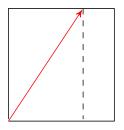
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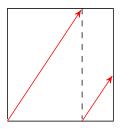
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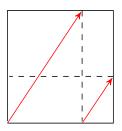
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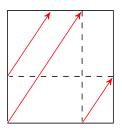
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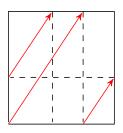
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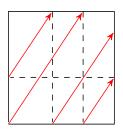
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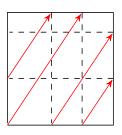
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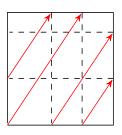
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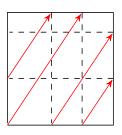
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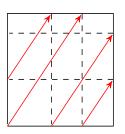
$$f(t) = 2 + \cos(t + \varphi_1) + \cos(\sqrt{2}t + \varphi_2) - e^{-t}$$



Consider the exponential polynomial

$$f(t) = \frac{2}{2} + \cos(t + \varphi_1) + \cos(\sqrt{2}t + \varphi_2) - e^{-t}$$

Kronecker's Theorem:  $(t, \sqrt{2}t) \mod 2\pi$  is dense in  $[0, 2\pi]^2$ 



Baker's Theorem:

$$\left|\left|\left(t,\sqrt{2}t\right)-\left(\pi-\varphi_1,\pi-\varphi_2\right)\right|\right|\geq \frac{1}{\mathrm{poly}(t)}$$

# Conclusion and Perspectives

#### The Discrete Case

A linear recurrence sequence is a sequence  $\langle u_0, u_1, u_2, \ldots \rangle$  of integers such that there exist constants  $a_1, \ldots, a_k$ , such that

$$u_{n+k} = a_1 u_{n+k-1} + a_2 u_{n+k-2} + \ldots + a_k u_n$$

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#### Theorem (Skolem 1934; Mahler 1935, 1956; Lech 1953)

The set of zeros of a linear recurrence sequence is semi-linear:

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#### Theorem (Berstel and Mignotte 1976)

In Skolem-Mahler-Lech, the infinite part (arithmetic progressions  $A_1, \ldots, A_\ell$ ) is fully constructive.

# The Skolem Problem

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"...a mathematical embarrassment ..."

Richard Lipton

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- Formidable "mathematical obstacle" at dimension 9 in the unbounded case.
- The infinite-zeros problem is also hard.
- Diophantine-approximation techniques unavoidable.